# Embedded PLC Hard Logic Solver Instruction Sets 

User's Manual

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## FOREWORD

Embedded PLC series bring the high performance, application flexibility and hardware compatibility to the Embedded PLC family of products. Ten contacts and abundant function blocks (also referred to as instructions) are provided for application control programs using the Embedded PLC series. In this manual, the usage for contacts and function blocks is described together with application examples.

## Contact elements include:

(1) $\dashv \vdash$ (A normally open contact, usually referred to as: "A contact")
(2) $\dashv / \vdash$ (A normally closed contact, usually referred to as: "B contact")
(3) -( )-(A normal coil)
(4) -(S)- (A set coil)
(5) -(R)- (A reset coil)
(6) -( $\uparrow$ )- (A positive transitional coil)
(7) -( $\downarrow$ )- (A negative transitional coil)
(8) -(M)- (A holding coil during power loss)
(9) -(SM)- (A holding set coil during power loss)
(10) -(RM)- (A holding reset coil during power loss)

## Function blocks instructions include:

(1) Timers and counters:

Timers: T1.0, T0.1, T0.01, Counters: UCTR, DCTR.
(2) Mathematical blocks:

Adders: ADD, ADDB, ADDL, ADBL, FADD, Subtracts: SUB, SUBB, SUBL, SBBL, FSUB, Multipliers: MUL, MULB, MULM, MUBM, MULL, MLBL, FMUL, Dividers: DIV, DIVB, DIVM, DVBM, DIVL, DVBL, FDIV, Square root: ISQR, FSQR.
(3) Register, Table, Array instructions:

Move: R->T, T->R, T->T, TXHG, BLKM, PACK,
Rotate/Shift: T_RS, BROT, ODSR,
Modify: MBIT,
Compare: T_CM, CMPR,
Search: T_SR,
Logic: AND, OR, XOR, COMP
Stack: PUSH, POP,
Sense: SENS,

Encoder, Decoder: ENCO, DECO, SSEG,
Convert: B->C, C->B, I->F, F->I.
(4) Flow control instructions:

Main program: EOP, SKIP, MCS, MSE, JMP, EOJ,
Subroutine: JSR, SBR, RET,
Loop: FOR, NEXT,
Pointer: INIP, INCP, DECP, PADD, PSUB.
(5) System related instructions:

System date: DGET, DSET, DCMP, System time: TGET, TSET, TCMP, System status: STAT.
(6) Others:

CAM
CDMR (CDM read)
CDMW (CDM write)
MOVE

Users are advised to become familiar with the binary operation (which can be found in any Digital Design Textbook) and the characteristics for each contact element and function block before designing a control application program. Please also be advised that the data and illustrations in this manual are not binding. We reserve the right to modify our products in line with our policy of continuous product improvement. Information in this manual is subject to change without notice and should not be treated as a commitment by FRECON Electric (Shenzhen) Co., Ltd. FRECON assumes no responsibility for any errors that may appear in this manual.

## CHAPTER 1: INTRODUCTION

The basic concept required to use this manual and the elements (contacts, function blocks, and instructions) in Embedded PLC is briefly described in this Chapter. In Section 1, the terminology and numerical representation are described. The constituents of a function block are described in Section 2 and the convention used to represent the function blocks is described in Section 3.

## SECTION 1: Terminology and Numerical Representations: BIT:

The basic unit of the binary system. The value of a bit is either 0 or 1 . The abbreviation for bit is B, such as $B 0, B 1, \ldots \ldots$. etc.

## NIBBLE:

A nibble is composed of four bits such as $\mathrm{B} 3 \sim \mathrm{~B} 0$. It can be used to represent decimal values ranging from 0 to 9 , or hexadecimal values ranging from $0 \sim \mathrm{~F}$. The abbreviation for nibble is NB , such as NB0, NB1, $\qquad$ etc.

## BYTE:

A byte is composed of eight bits (B7~B0) or two contiguous nibbles (NB1~NB0). It can be used to represent hexadecimal values ranging from $00 \sim \mathrm{FF}$. The abbreviation for byte is BY, such as BY0, BY1, $\qquad$ etc.

## WORD:

A word is composed of sixteen bits. It can be used to represent hexadecimal values ranging from $0000 \sim$ FFFF or $0 \sim 65535$ in the decimal system. The abbreviation for word is W, such as W0, W1, ..... etc. Since Embedded PLC is based on 16-bit microcomputer architecture, a word occupies one register in the computer memory.

## LONG WORD:

A long word is composed of two continuous words or 32 bits. It can be used to represent hexadecimal values ranging from 00000000~FFFFFFFF, floating point numbers through special convention, or decimal format ranging from 0~99999999. The abbreviation for long word is LW, such as LW0, LW1, $\qquad$ etc. A long word occupies two continuous registers in the computer memory. The first register contains the most significant 16 bits (usually referred to as HIGH WORD), the second register contains the least significant 16 bits (usually referred to as LOW WORD). A long word is referenced by the address occupied by the High Word.

## Floating Point Representation using a Long Word:

A long word ( 32 bits) can be used to represent a floating point number. The bit assignment is shown in the following figure:


Formula: $\quad \mathbf{I}=(\mathbf{- 1})^{\mathbf{S}} \times \mathbf{2}^{(\mathrm{E}-64)} \times \mathbf{F r}$

For example, assuming that the content of register 40130 is C 000 h and register 40131 is 0042 h ; then for an operation using floating point referencing register 40130 (40130 and 40131 actually), the value used is:

40130
40131
$1100000000000000 \quad 0000000001000010$
$I=(-1)^{0} \times 2^{(66-64)} \times\left(2^{-1}+2^{-2}\right)=3$

## SECTION 2: Constituents of a Function Block

In Embedded PLC series, a function block is composed of four parts: Function Name, Input Control, Operand and Function Output as shown in the following figure:


Where: 1. $\mathrm{I}_{1}, \mathrm{I}_{2}, \mathrm{I}_{3}$ are Input controls
2. $\mathrm{O}_{1}, \mathrm{O}_{2}, \mathrm{O}_{3}$ are Function outputs.
3. TOP, MIDDLE, BOTTOM stand for Top node, middle node and bottom node. These three nodes are operands.
4. NAME is the name of the function block.

## Function Name:

The function name is an abbreviation or acronym of the operation performed by the function block. Two to four characters are used to represent the function. A complete list of the function block names may be found in the FOREWORD of this manual.

## Input Control:

There must be one input control for each function block. This input control (usually referred to as $I_{1}$ ) is used to determine whether to execute this function block or not. For some function blocks, there are two additional input controls ( $\mathrm{I}_{2}$ and $\mathrm{I}_{3}$ ). They are used to determine the execution mode of the function block.

## Function Output:

There must be a function output control for each function block. This output (usually referred to as $\mathrm{O}_{1}$ ) is used to drive a coil or used as an input control for the next function block. For some function blocks, there are two additional output controls $\left(\mathrm{O}_{2}\right.$ and $\left.\mathrm{O}_{3}\right)$, they are also used to represent the results of the execution.

## Operands:

Operands, as the name implies, are the objects of operations. An operand whose content is not altered by the operation is called a SOURCE. An operand that is used to store the result of the operation is called a DESTINATION. Operands can be Input contact, Output coil or register in memory. For Embedded PLC, the designations of operands are listed in the following table:

Table 1.1: Operands

| Initial | NAME | DESCRIPTION |
| :---: | :---: | :--- |
| $\mathbf{0}$ | Output Coil <br> (Discrete output) | Use Output coil as an operand. Since 1 word $=16$ bits, thus the <br> number assignment of the operand must be a multiple of 16 plus 1. <br> For example: 00001, 00017, 00033. |
| $\mathbf{1}$ | Input contact <br> (Discrete input) | Use Input contact as an operand. The number assignment of the <br> operand must be a multiple of 16 plus 1. For example: 10001, <br> $10017,10033$. |
| $\mathbf{3}$ | Input register | Use Input register as an operand. For example:30001, 30003. |
| $\mathbf{4}$ | Holding register | Use Holding register as an operand. For example:40001, 40003. |
| $\mathbf{C}$ | Constant | For some function blocks, a constant can be defined as an operand: <br> and during control program execution, the value of the constant is <br> readily available rather then fetching from register memory. For |
| $\mathbf{P}$ | Pointer | example: \#00001, \#0020h. The former is a decimal constant, and <br> the latter is a hexadecimal constant. |
| For some function blocks, a pointer can be defined as an operand, |  |  |
| and this pointer can be used for indirect addressing pointing to 0-, |  |  |
| $1-, 3-, 4-t y p e ~ v a r i a b l e . ~ F o r ~ e x a m p l e: ~ P 0001 ~$ |  |  |, | For paired instructions (such as FOR and NEXT), their operands |
| :--- |
| are label, and the label for each instruction must be the same in |
| order for program to be executed correctly. For example: L0001. |,

Currently, there are three models of Embedded-PLC controllers. The memory size and the CPU capability are different between models to meet different control requirements. Therefore, the numbers of spaces available for operands are also different. The available ranges for operands for each model are listed in the following table.

Table 1.2: Available operand ranges for different models of the Embedded PLC series controller

| OPERAND | Embedded PLC |
| :---: | :---: |
| $\mathbf{0}$ | $00001 \sim 09984$ |
| $\mathbf{1}$ | $10001 \sim 12048$ |
| $\mathbf{3}$ | $30001 \sim 30512$ |
| $\mathbf{4}$ | $40001 \sim 49999$ |
| $\mathbf{L}$ | L1~L150 |
| $\mathbf{C}$ | $0 \sim 65535$ |
| $\mathbf{P}$ | P0~P15 |

## CHAPTER 2：CONTACTS

Contact elements are the most fundamental elements in Ladder Programs．Familiarization with their characteristics and usage is highly recommended．

## （1）$\dashv \vdash$ Normally Open Contact：

This type of contact is usually referred to as＂A Contact＂．When a contact is energized，the said＂A contact＂becomes conductive；and vice versa．


## 【Meaning】

When input contact 10001 is＇ON＇，coil 00001 is energized，and＂A contact＂ 00001 becomes conductive，thus，coil 00002 is energized．

## （2）$\dashv / \vdash$ Normally Closed Contact：

This type of contact is usually referred to as＂B Contact＂．When a contact is not energized，the said ＂B contact＂becomes conductive；and vice versa．


## 【Meaning】

When input contact 10001 is＇OFF＇，coil 00001 is energized，and＂B contact＂ 00001 becomes non－conductive，thus，coil 00002 is not energized．
（3）－（ ）－Output Coil：
This output coil reflects the state of the elements connected to it．If the element is in the＇ ON ＇state， then this coil is said to be energized；and vice versa．

【EXAMPLE】

| $\mid r$ |  |
| ---: | ---: |
| 10001 | 00001 |

【Meaning】When input contact 10001 is＇ON＇，then output coil 00001 is＇ ON ＇；When input contact 10001 is＇OFF＇，then output coil 00001 is＇OFF＇

【Timing diagram】

（4）－（S）－Set Coil ：
When the element connected to this coil is＇ON＇，then this set coil is set to＇ON＇and remains in that ＂ON＇state until the＂RESET coil＂with the same reference number is energized．


## 【Meaning】

When contact 10001 is＇ ON ＇，the set coil 00001 is＇ ON ＇and remains＇ON＇no matter how contact 10001 is changed．

（5）－（R）－Reset coil：
When the element connected to this coil is＇ON＇，then this set coil is set to＇OFF＇and remains in that ＂OFF＇state until the＂SET coil＂with the same reference number is energized．


## 【Meaning】

When input contact 10001 is＇ON＇，output coil 00001 is set to＇ON＇and remains in that state．Until input contact 10001 is＇OFF＇and input contact 10002 is＇ ON ＇，then output coil 00001 is set to＇OFF＇ and remains＇OFF＇．

## 【Timing diagram】


（6）－（ $\uparrow$ ）－Positive Transitional Pulse Output Coil：
When the element connected to this output has an＇OFF $\square^{\prime}$ ON＇transition，a pulse（＇OFF＇$\square$＇ON＇）is generated for this output．


## 【Meaning】

When input contact 1000 receives a transition＇ $\mathrm{OFF}^{\prime} \square$＇ $\mathrm{ON}^{\prime}$ ，then a pulse＇ $\mathrm{OFF}{ }^{\prime} \square$＇ $\mathrm{ON}^{\prime}$＇is generated for output coil 00001 ．The width of the pulse is 1 scan time．

## 【Timing diagram】


（7）－（ $\downarrow$ ）－Negative Transitional Pulse Output Coil：
When the element connected to this output has an＇ON＇$\square$＇OFF＇transition，a pulse（＇OFF $\square$＇ON＇）is generated for this output．


## 【Meaning】

When input contact 10001 receives a transition＇ON＇$\square$＇OFF＇，then a pulse＇OFF＇$\square$＇ON＇is generated for output coil 00001 ．The width of the pulse is 1 scan time．

【Timing diagram】

(8) -(M)- Holding Coil during power loss:

This output coil reflects the state of the elements connected to it. If the element is in the 'ON' state, then this coil is said to be energized; and vice versa. The last state of the coil is maintained after system power is shut down and turned on again.
(9) -(SM)- Holding Set Coil during power loss:

When the element connected to this coil is 'ON', then this coil is set to ' ON ' and remains in that "ON' state until the "RESET coil" with the same reference number is energized. The last state of the coil is maintained after system power is shut down and turned on again.
(10) -(RM)- Holding Reset Coil during power loss:

When the element connected to this coil is 'ON', then this coil is set to 'OFF' and remains in that "OFF' state until the "SET coil" with the same reference number is energized. The last state of the coil is maintained after system power is shut down and turned on again.

## CHAPTER 3: FUNCTION BLOCKS



Description:


Node description:
TOP :


MIDDLE :
BOTTOM :
Input Control:
$\mathrm{I}_{1}:$
$\mathrm{I}_{2}:$
$\mathrm{I}_{3}:$
Function Output:
$\mathrm{O}_{1}$ : $\qquad$
$\mathrm{O}_{2}$ :
$\square$
$\mathrm{O}_{3}$ :

The template for the description of a function block is divided into ten areas ( $\square \sim \square$ ). The meaning for each area is described as follows:
$\square$ NAME:
NAME is an abbreviation or acronym for the operation performed by the function block. Two to four characters are used to represent the function. When displaying the ladder program on screen, the name of the function block is also displayed.
$\square$ Full Name of Function Block:
The operation of the function block is given briefly in this area.

- NAME:

This area is provided for easy reference to function blocks.

## $\square$ Trigger mode:

The entry here is used to indicate the trigger mode of the function block. For "Level trigger" mode, when I1 is HIGH, then the function block is executed. For "Edge trigger" mode, when there is an OFF to ON transition, then the function block is executed. For edge-trigger function blocks, a " ^" mark is prefixed to the name of the function block in the PP programming environment.

## $\square$ Symbol:

The symbol of the function block as used in this manual is displayed in the ladder diagram.

## $\square$ Operands

Operands available for the function block have a circle "O " marked in the table.
$\square$ Function block description:
A brief description of the major function of the function block together with its input control, function output and result of the execution is given in this area.
$\square$ Node description:
The usage of each node, whether it is a Source or a Destination, is given in this area.
$\square$ Input Control:
The condition $\left(I_{1}\right)$ required for the function block to be executed is described here. The execution mode ( $\mathrm{I}_{2}$ and/or $\mathrm{I}_{3}$ ) is also described here.
$\square$ Function Output:
The results of the execution $\left(\mathrm{O}_{1,} \mathrm{O}_{2}, \mathrm{O}_{3}\right)$ are given in this area.

| T1.0 | 1.0 SECOND TIMER | $\curvearrowleft$ |  |
| :--- | :--- | :--- | :--- |

## SYMBOL:



OPERANDS:

(1)0~65535

## Description:

Timer increments by one at intervals of one second. When the accumulated time (stored in the BOTTOM node) reaches the timer preset (stored in TOP node), the timer stops. Input control can be used to start, stop and reset the timer. The timer status (whether the elapsed time has reached the preset time) can be detected by examining the function output.

## Node description:

TOP: Preset value for timer.
BOTTOM: Accumulated value since timer started.

## Input Control:

$\mathrm{I}_{1}$ : Execution control. When $\mathrm{I}_{1}=1$, timer starts; $\mathrm{I}_{1}=0$, timer stops.
$I_{2}$ : Reset control, when $I_{2}=0$, the accumulated value is cleared to zero.

## Function Output:

$\mathrm{O}_{1}=1$, if accumulated value $=$ preset value.
$=0$, if accumulated value $<$ preset value.
$\mathrm{O}_{2}$ : Complement of $\mathrm{O}_{1}$

## 【EXAMPLE】



## 【DESCRIPTION】

This example shows a five－second timer．The decomposition of actions are：
1． 40012 is 0 ，then $00040=$＇OFF＇and $00041=$＇ON＇at the beginning．
2．When input control 10012 is＇ ON ＇，register 40012 increments by one for every one second．
3．When the content of register $40012=5$（as defined in the top node），the function output：
$00040=$＇ON＇， $00041=$＇OFF＇．
4．Since $00040=$＇ ON ＇， $\mathrm{I}_{2}$ changes to＇OFF＇，and clears register 40012 to＇ 0 ＇．
5．Since $40012=0$ ，then $00040=$＇OFF＇， $00041=$＇ON＇，register 40012 continues incrementing， and the execution continues from STEP 3.

| T0.1 | 0.1 SECOND TIMER | $\checkmark\llcorner$ |  |
| :---: | :---: | :---: | :---: |

SYMBOL:


OPERANDS:

(1) $0 \sim 65535$

## Description:

Timer increments by one at intervals of 0.1 second. When the accumulated time (stored in the BOTTOM node) reaches the timer preset (stored in TOP node), the timer stops. Input control can be used to start, stop and reset the timer. The timer status (whether the elapsed time has reached the preset time) can be detected by examining the function output.

## Node description:

TOP: Preset value for timer.
BOTTOM: Accumulated value since timer started.

## Input Control:

$\mathrm{I}_{1}$ : Execution control. When $\mathrm{I}_{1}=1$, timer starts; $\mathrm{I}_{1}=0$, timer stops.
$\mathrm{I}_{2}$ : Reset control, when $\mathrm{I}_{2}=0$, the accumulated value is cleared to zero.

## Function Output:

$\mathrm{O}_{1}=1$, if accumulated value $=$ preset value.
$=0$, if accumulated value $<$ preset value.
$\mathrm{O}_{2}$ : Complement of $\mathrm{O}_{1}$

## 【EXAMPLE】



## 【DESCRIPTION】

This example shows a five－second timer．The decomposition of actions are：

1． 40012 is 0 ，then $00040={ }^{\prime}$＇OFF＇and $00041=' \mathrm{ON}$＇at the beginning．
2．When input control 10012 is＇ ON ＇，register 40012 increments by one for every one second．
3．When the content of register $40012=50$（as defined in the top node），the function output：
$00040=$＇ON＇， $00041=$＇OFF＇．
4．Since $00040=$＇ ON ＇， $\mathrm{I}_{2}$ changes to＇OFF＇，and clears register 40012 to＇ 0 ＇．
5．Since $40012=0$ ，then $00040=$＇OFF＇， $00041=$＇ON＇，register 40012 continues incrementing， and the execution continues from STEP 3.

| T0.01 | 0.01 SECOND TIMER | $\checkmark\llcorner$ |  |
| :---: | :---: | :---: | :---: |

SYMBOL:


OPERANDS:

(1)0~65535

## Description:

Timer increments by one at intervals of 0.01 second. When the accumulated time (stored in the BOTTOM node) reaches the timer preset (stored in TOP node), the timer stops. Input control can be used to start, stop and reset the timer. The timer status (whether the elapsed time has reached the preset time) can be detected by examining the function output.

## Node description:

TOP: Preset value for timer.
BOTTOM: Accumulated value since timer started.

## Input Control:

$I_{1}$ : Execution control. When $I_{1}=1$, timer starts; $I_{1}=0$, timer stops.
$I_{2}:$ Reset control, when $I_{2}=0$, the accumulated value is cleared to zero.

## Function Output:

$\mathrm{O}_{1}=1$, if accumulated value $=$ preset value.
$=0$, if accumulated value $<$ preset value .
$\mathrm{O}_{2}$ : Complement of $\mathrm{O}_{1}$

## 【EXAMPLE】



## 【DESCRIPTION】

This example shows a five－second timer．The decomposition of actions are：

1． 40012 is 0 ，then $00040=$＇OFF＇and $00041=$＇ON＇at the beginning．
2．When input control 10012 is＇ ON ＇，register 40012 increments by one for every 0.01 second．
3．When the content of register $40012=500$（as defined in the top node），the function output：
$00040=$＇ON＇， $00041=$＇OFF＇．
4．Since $00040=$＇ ON ＇， $\mathrm{I}_{2}$ changes to＇OFF＇，and clears register 40012 to＇ 0 ＇．
5．Since $40012=0$ ，then $00040=$＇OFF＇，register 40012 continues incrementing，and the execution continues from STEP 3.

| UCTR | UP COUNTER | $f /$ |
| :--- | :--- | :--- | :--- |

SYMBOL:


OPERANDS:

(1)0~65535

## Description:

This counter counts the pulses presented at $\mathrm{I}_{1}$ from 0 to a preset value. Input control can be used to start, stop and reset the counter. The counter status (whether the accumulated value has reached the preset value) can be detected by examining the function output.

## Node description:

TOP: Preset value for counter.
BOTTOM: Accumulated value since counter started.

## Input Control:

$\mathrm{I}_{1}$ : Counter control. When $\mathrm{I}_{1}$ receives an 'OFF' $\square$ 'ON' transition, The counter is incremented by 1.
$\mathrm{I}_{2}$ : Reset control. When $\mathrm{I}_{2}=0$, the accumulated value is cleared to zero.

## Function Output:

$\mathrm{O}_{1}=1$, if accumulated value=preset value.
$=0$, if accumulated value < preset value.
$\mathrm{O}_{2}$ : Complement of $\mathrm{O}_{1}$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 receives an OFF to ON transition，the accumulated value of the counter （40001）is incremented by 1 ．When the accumulated value reaches 100 ，coil 00001 is energized． When normal－close contact 00001 opens，the counter is reset．

SYMBOL:


OPERANDS:

(1)0~65535

## Description:

This counter counts the pulses presented at $\mathrm{I}_{1}$ from a preset value to 0 . Input control can be used to start, stop and reset the counter. The counter status (whether the accumulated value has reached 0 ) can be detected by examining the function output.

## Node description:

TOP: Preset value for counter.
BOTTOM: Accumulated value since counter started.

## Input Control:

$\mathrm{I}_{1}$ : Counter control. When $\mathrm{I}_{1}$ receives an 'OFF' $\square^{\prime} \mathrm{ON}^{\prime}$ transition, the counter is decreased by 1.
$\mathrm{I}_{2}$ : Reset control. When $\mathrm{I}_{2}=0$, the accumulated value is set to preset value.

## Function Output:

$\mathrm{O}_{1}=1$, if accumulated value $=0$.
$=0$, if accumulated value $>0$.
$\mathrm{O}_{2}$ : Complement of $\mathrm{O}_{1}$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 receives an OFF to ON transition，the accumulated value of the counter （40001）is decreased by 1 ．When the accumulated value reaches 0 ，coil 00001 is energized．When normal－close contact 00001opens，the counter is reset，and the value in counter（40001）is set to 100.

|  |  |  |  |
| :---: | :---: | :---: | :---: |
| ADD | FOUR DIGIT DECIMAL ADDER | $\boxed{ } 1$ | $\uparrow \square$ |

## SYMBOL:



OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  | $\bigcirc$ | $\bigcirc$ | $(1$ | $\bigcirc$ |  |
| MIDDLE |  |  | $\bigcirc$ | $\bigcirc$ | $(1$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1)0~9999

## Description:

The decimal values stored in the top and middle nodes are added and the sum is stored in the bottom node.
Sum $=\left(\right.$ top + middle $\left.+\mathrm{I}_{3}\right)$ MOD 10000.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function output $\left(\mathrm{O}_{3}\right)$ may be used to determine whether or not an overflow has occurred.

## Node Description:

TOP: Summand, must be $<10000$.
MIDDLE: Addend, must be $<10000$.
BOTTOM:1.(top + middle $+\mathrm{I}_{3}$ ) MOD 10000
2.If error (ref. to $\mathrm{O}_{2}$ ) occurred, the content of the bottom node remains unchanged.

## Input Control:

$\mathrm{I}_{1}$ : When $\sqrt{ }$ ( $\uparrow)$ is presented, the function block is executed.
$\mathrm{I}_{2}$ : error in
$\mathrm{I}_{3}$ : carry in

## Function Output:

$$
\mathrm{O}_{1}=\mathrm{I}_{1}
$$

$\mathrm{O}_{2}$ = error output ( $\mathrm{O}_{2}$ is ' 1 ' if $\mathrm{I}_{2}$ is ' 1 ' or the value of either top node or middle node is over 9999)
$\mathrm{O}_{3}$ : overflow/carry
$=1$, Sum $>9999$
$=0$, Sum $\leqq 9999$

## 【EXAMPLE】



## 【DESCRIPTION】

When the contact 10025 has an＂OFF $\rightarrow$ ON＂，the content of register 40021 is added to the content of register 40027 and the sum is stored back to register（40021）．Since the sum is larger than 9999 ，therefore，the second adder is energized．


## ADDB

FOUR DIGIT HEXADECIMAL ADDER

SYMBOL:


OPERANDS

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  | $\bigcirc$ | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| MIDDLE |  |  | $\bigcirc$ | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1) $0 \sim 65535$
word + word $\rightarrow$ word (hexadecimal)

## Description:

The values (hexadecimal) stored in the top and middle nodes are added and the sum is stored in the bottom node. Sum = (top + middle) MOD 65536.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function output $\left(\mathrm{O}_{3}\right)$ may be used to determine whether or not an overflow has occurred.

## Node Description:

TOP: Summand, must be $<65535$.
MIDDLE: Addend, must be $<65535$.
BOTTOM: Sum < 65535. The carry, if any, is ignored.

## Input Control:

$\mathrm{I}_{1}$ : When $\rfloor(\jmath \square)$ is presented, the function block is executed.

## Function Output:

$$
\begin{aligned}
\mathrm{O}_{1} & =\mathrm{I}_{1} \\
\mathrm{O}_{2} & =0 \\
\mathrm{O}_{3} & : \text { overflow } \\
& =1, \text { Sum }>65535 \\
& =0, \text { Sum } \leqq 65535
\end{aligned}
$$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10025 is energized，the content of register 40021 is added to the content of register 40027 and the sum is stored back to register 40021．Since the sum is larger than 65535 ， therefore，the second adder is energized．


## ADDL

EIGHT DIGIT DECIMAL ADDER

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE |  |  |  | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1) 0~9999

$$
\text { Lword }+ \text { Lword } \rightarrow \text { Lword (Decimal) }
$$

## Description:

The values (long word, decimal) stored in the top and middle nodes are added and the sum is stored in the bottom node. $\mathrm{Sum}=\mathrm{MOD}\left(\right.$ top + middle $\left.+\mathrm{I}_{3}\right) \mathrm{MOD} 100000000$.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function output $\left(\mathrm{O}_{3}\right)$ may be used to determine whether or not an overflow has occurred.

## Node Description:

TOP: Summand, must be $<100000000$.
MIDDLE: Addend, must be $<100000000$.
BOTTOM :1. (top + middle $+\mathrm{I}_{3}$ ) MOD 10000000.
2.If error (refer to $\mathrm{O}_{2}$ ) occurred, the content of the bottom node remains unchanged.

## Input Control:

$\mathrm{I}_{1}$ : When $\sqrt{ }$ ( $\left.\uparrow L\right)$ is presented, the function block is executed.
$\mathrm{I}_{2}$ : error in
$\mathrm{I}_{3}$ : carry in

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}$ = error output ( $\mathrm{O}_{2}$ is ' 1 ' if $\mathrm{I}_{2}$ is ' 1 ' or the value of either top node or middle node is over 99999999.)
$\mathrm{O}_{3}$ : overflow/carry
$=1$, Sum $>99999999$
= 0, Sum $\leqq 99999999$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10025 is＂ON＇，the content of registers 40027／40028 is added to the content of registers 40021／40022．The sum is stored in registers 40030／40031．Since the sum is less then 99999999，thus， $\mathrm{O}_{1}: \mathrm{ON}, \mathrm{O}_{2}=\mathrm{O}_{3}=\mathrm{OFF}$ ．


| ADBL | EIGHT DIGIT HEXADECIMAL ADDER | $\checkmark L$ | $\uparrow R$ |
| :--- | :--- | :--- | :--- |

SYMBOL:


OPERANDS

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE |  |  |  | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1)0~65535

Lword + Lword $\rightarrow$ Lword (Hexadecimal)

## Description:

The values (long word, hexadecimal) stored in the top and middle nodes are added and the sum is stored in the bottom node. Sum = MOD(top+middle)FFFFFFFF.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function output $\left(\mathrm{O}_{3}\right)$ may be used to determine whether or not an overflow has occurred.

## Node Description:

TOP: Summand, must be $<100000000_{\text {hex }}$.
MIDDLE: Addend, must be $<100000000_{\text {hex }}$.

Input Control:
$\mathrm{I}_{1}$ : When $\Omega(\uparrow L)$ is presented, the function block is executed.

## Function Output:

$$
\begin{aligned}
\mathrm{O}_{1} & =\mathrm{I}_{1} \\
\mathrm{O}_{2} & =0 \\
\mathrm{O}_{3} & : \text { overflow } \\
& =1, \text { Sum }>4294967295\left(=100000000_{\text {hex }}\right) \\
& =0, \text { Sum } \leqq 4294967295\left(=100000000_{\text {hex }}\right)
\end{aligned}
$$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10025 is＂ ON ＇，the content of registers $40027 \& 40028$ is added to the content of registers $40021 / 40022$ ．The sum is stored in the registers $40030 \& 40031$.


| FADD | FLOATING POINT ADDER | $\lrcorner \square$ | $\uparrow \mathrm{L}$ |
| :---: | :---: | :--- | :--- |

## SYMBOL:



OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE |  |  |  | $\bigcirc$ | $\square$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

$\square 0 \sim 65535$
float + float $\rightarrow$ float

## Description:

The values (floating point) stored in the top and middle nodes are added and the sum is stored in the bottom node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.

## Node Description:

TOP: Summand.
MIDDLE: Addend.
BOTTOM: Sum.

## Input Control:

$\mathrm{I}_{1}$ : When $\qquad$ ( $\uparrow \square$ ) is presented, the function block is executed.

## Function Output:

$$
\begin{aligned}
& \mathrm{O}_{1}=\mathrm{I}_{1} \\
& \mathrm{O}_{2}=0 \\
& \mathrm{O}_{3}=0
\end{aligned}
$$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10025 is＂ON＇，the content of registers 40010／40011 is added to the content of registers 40020／40021；the sum is stored in registers 40030／40031；and $\mathrm{O}_{1}: \mathrm{ON}, \mathrm{O}_{2}=\mathrm{O}_{3}=\mathrm{OFF}$ ．


| SUB | FOUR DIGIT DECIMAL SUBTRACTOR | $\checkmark \square$ | $\uparrow \mathrm{L}$ |
| :---: | :---: | :---: | :---: |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  | $\bigcirc$ | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| MIDDLE |  |  | $\bigcirc$ | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1)0~9999

$$
\text { word }- \text { word } \rightarrow \text { word (Decimal) }
$$

## Description:

The value stored in the middle node is subtracted from the top node, and the difference is stored in the bottom node.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function output may be used to determine the relationship between minuend and subtrahend ( $>,=,<$ ).

## Node Description:

TOP: Minuend, must be $<10000$.
MIDDLE: Subtrahend, must be $<10000$.
BOTTOM: Difference.

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark(\uparrow \square)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=1$, if difference $>0$ (Top node $>$ Middle node).
$\mathrm{O}_{2}=1$, if difference $=0($ Top node $=$ Middle node $)$.
$\mathrm{O}_{3}=1$, if difference $<0$ (Top node $<$ Middle node).

## 【EXAMPLE】



## 【DESCRIPTION】

Assume that register（40001）$=9000_{(10)}$ ，and（40002）$=500_{(10)}$ ．when contact 00080 is＇ON＇，the subtraction： $\mathbf{( 4 0 0 0 3 )}=(\mathbf{4 0 0 0 1})-(\mathbf{4 0 0 0 2})$ is performed．Since the minuend is larger than the subtrahend，thus coil 00011 is＇ON＇， 00012 is＇OFF＇and 00013 is＇OFF＇．
$\left.\begin{array}{l|l|l}40001 \\ 40002 \\ 4003\end{array}\right)$

| SUBB | FOUR DIGIT HEXADECIMAL <br> SUBTRACTOR | $\checkmark \square$ | $\uparrow L$ |
| :---: | :---: | :---: | :---: |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  | $\bigcirc$ | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| MIDDLE |  |  | $\bigcirc$ | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1)0~65535
word - word $\rightarrow$ word (Binary)

## Description:

The value stored in the middle node is subtracted from the top node, and the difference is stored in the bottom node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function output may be used to determine the relationship between minuend and subtrahend ( $>$, = , $<$ ).

## Node Description:

TOP: Minuend
MIDDLE: Subtrahend
BOTTOM: Difference.

## Input Control:

$\mathrm{I}_{1}$ : When $\Omega(\uparrow L)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=1$, if difference $>0$.
$\mathrm{O}_{2}=1$, if difference $=0$.
$\mathrm{O}_{3}=1$, if difference $<0$.

## 【EXAMPLE】



## 【DESCRIPTION】

Assume that register $(40001)=9000_{(10)}$ ，and（40002）$=9000_{(10)}$ ．when contact 00080 is＇ ON ＇，the subtraction：$(\mathbf{4 0 0 0 3})=(\mathbf{4 0 0 0 1})-(\mathbf{4 0 0 0 2})$ is performed．Since the minuend is equal to the subtrahend，thus coil 00012 is＇ON＇．

| 40001 | 09000 | 9000 |
| :---: | :---: | :---: |
| 40002 | 09000 | － 9000 |
| 40003 | 00000 | 0000 |


| SUBL | EIGHT DIGIT DECIMAL SUBTRACTOR | $\checkmark L$ | $\uparrow L$ |
| :---: | :---: | :--- | :--- |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  |  | $\bigcirc$ | $\square$ | $\bigcirc$ |  |
| MIDDLE |  |  |  | $\bigcirc$ | $\square$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

$\square 0 \sim 65535$

## Lword-Lword $\rightarrow$ Lword (Decimal)

## Description:

The value stored in the middle node is subtracted from the top node, and the difference is stored in the bottom node.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function output may be used to determine the relationship between minuend and subtrahend ( $>,=,<$ ).

## Node Description:

TOP: Minuend, must be <=99999999.
MIDDLE: Subtrahend, must be $<=99999999$.
BOTTOM: Difference.

Input Control:
$\mathrm{I}_{1}$ : When $\square(\uparrow L)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=1$, if difference $>0$.
$\mathrm{O}_{2}=1$, if difference $=0$.
$\mathrm{O}_{3}=1$, if difference $<0$.

## 【EXAMPLE】



## 【DESCRIPTION】

Assume that long $\operatorname{word}(40010)=9999_{(10)}$ and long word $(40020)=9999_{(10 .)}$ ．when contact 00080 is ＇ON＇，the operation：long word（40030）＝long word（40010）－long word $(\mathbf{4 0 0 2 0})$ is performed． Since the minuend is equal to the subtrahend，thus coil 00012 is＇ ON ＇．

| 40010 | 0000 |
| :--- | :--- |
| 40011 | 9999 |
|  |  |
| 40020 | 0000 |
| 40021 | 9999 |
|  |  |
| 40030 | 0000 |
| 40031 | 0000 |
|  |  |


| SBBL | EIGHT DIGIT HEXADECIMAL <br> SUBTRACTOR | $\checkmark\llcorner$ | $\uparrow\llcorner$ |
| :---: | :---: | :---: | :---: |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  |  | $\bigcirc$ | $\square$ | $\bigcirc$ |  |
| MIDDLE |  |  |  | $\bigcirc$ | $\square$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

$\square 0 \sim 65535$

$$
\text { Lword - Lword } \rightarrow \text { Lword(Binary) }
$$

## Description:

The value stored in the middle node is subtracted from the top node, and the difference is stored in the bottom node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function output may be used to determine the relationship between minuend and subtrahend ( $>,=,<$ ).

## Node Description:

TOP: Minuend.
MIDDLE: Subtrahend.
BOTTOM: Difference .

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark L(\uparrow L)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=1$, if difference $>0$.
$\mathrm{O}_{2}=1$, if difference $=0$.
$\mathrm{O}_{3}=1$, if difference $<0$.

## 【EXAMPLE】



## 【DESCRIPTION】

Assume that long $\operatorname{word}(40010)=65536(10)$ and $\operatorname{long} \operatorname{word}(40020)=65536(10)$ ．when contact 00080 is＇ ON ＇，the operation：long word（40030）＝long word（40010）－long word $(\mathbf{4 0 0 2 0})$ is performed． Since the minuend is equal to the subtrahend，thus coil 00012 is＇ ON ＇．

| 40010 | 0001 |
| :--- | :--- |
| 40011 | 0000 |
|  |  |
| 40020 | 0001 |
| 40021 | 0000 |
|  |  |
| 40030 | 0000 |
| 40031 | 0000 |
|  |  |

FSUB
FLOATING POINT SUBTRACTOR

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  |  | $\bigcirc$ | $\square$ | $\bigcirc$ |  |
| MIDDLE |  |  |  | $\bigcirc$ | $\square$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

$\square 0 \sim 65535$

Float - float $\rightarrow$ float

## Description:

The value stored in the middle node is subtracted from the top node, and the difference is stored in the bottom node.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function output may be used to determine the relationship between minuend and subtrahend ( $>,=,<$ ).

## Node Description:

TOP: Minuend.
MIDDLE: Subtrahend.
BOTTOM: Difference.

## Input Control:

$\mathrm{I}_{1}$ : When $\lceil(\uparrow \square)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=1$, if difference $>0$.
$\mathrm{O}_{2}=1$, if difference $=0$.
$\mathrm{O}_{3}=1$, if difference $<0$.

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10025 is＂ON＇，the content of registers 40020／40021 is subtracted from the content of registers 40010／40011；the difference is stored in registers 40030／40031．Since the minuend is greater than the subtrahend，thus $\mathrm{O}_{1}: \mathrm{ON}, \mathrm{O}_{2}=\mathrm{O}_{3}=\mathrm{OFF}$ ．


|  |  | MUL |  |
| :---: | :---: | :---: | :---: |
| MUL | FOUR DIGIT DECIMAL MULTIPLIER | $\boxed{ }$ | $\uparrow \square$ |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  | $\bigcirc$ | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE |  |  | $\bigcirc$ | $\bigcirc$ | $(1$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1)0~9999
word $\times$ word $\rightarrow$ Lword (Decimal)

## Description:

The value in the top node is multiplied by the value in the middle node, and the product is stored in the bottom node(long word).
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function output $\left(\mathrm{O}_{3}\right)$ may be used to determine whether or not an overflow has occurred.

## Node Description:

TOP: Multiplicand, must be $<=9999$.
MIDDLE: Multiplier, must be <=9999.
BOTTOM :1.Product, Long word.
2.If error (refer to $\mathrm{O}_{2}$ ) occurred, the content of the bottom node remains unchanged.

## Input Control:

$\mathrm{I}_{1}$ : When $\sqrt{ }$ $(\uparrow \quad)$ is presented, the function block is executed.
$\mathrm{I}_{2}$ : error in

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=1$ (If the value of either top node or middle node is greater than 9999).
$\mathrm{O}_{3}: 0$

## 【EXAMPLE】



## 【DESCRIPTION】

Let register（40005）＝2500 and（40006）$=1100$ ．When contact 10007 is＇ON＇，the operation： long word $(40036)=(40005) \times(40006)$ is performed．


SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  | $\bigcirc$ | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE |  |  | $\bigcirc$ | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1)0~65535
word $\times$ word $\rightarrow$ Lword (Binary)

## Description:

The value in the top node is multiplied by the value in the middle node, and the product is stored in the bottom node(long word).
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.

## Node Description:

TOP: Multiplicand.
MIDDLE: Multiplier.
BOTTOM: Product, Long word.
Input Control:
$\mathrm{I}_{1}$ : When $\checkmark L(\uparrow L)$ is presented, the function block is executed.

## Function Output:

$$
\begin{aligned}
& \mathrm{O}_{1}=\mathrm{I}_{1} \\
& \mathrm{O}_{2}=0 \\
& \mathrm{O}_{3}=0
\end{aligned}
$$

## 【EXAMPLE】



## 【DESCRIPTION】

Let register $(40005)=2500$ and $(40006)=1100$ ．When contact 10007 is＇$O N$＇，the operation： long word $(40036)=(40005) \times(40006)$ is performed．

| 40005 | 2500 | $2500 \times 1100=2750000$ |
| :---: | :---: | :---: |
| 40006 | 1100 | The product is stored in registers 40036／40037． |
| 40036 | $0029{ }_{\text {（hex }}$ |  |
| 40037 | F630 ${ }_{\text {（hex）}}$ |  |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  | $\bigcirc$ | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE |  |  | $\bigcirc$ | $\bigcirc$ | $\bigcirc$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1)0~9999
word $\times$ word $\rightarrow$ word (Decimal)

## Description:

The value in the top node is multiplied by the value in the middle node, and the product is stored in the bottom node(long word).

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function output $\left(\mathrm{O}_{3}\right)$ may be used to determine whether or not an overflow has occurred .

## Node Description:

TOP: Multiplicand (<= 9999).
MIDDLE: Multiplier(<=9999).
BOTTOM:1.Product.
2.If error (refer to $\mathrm{O}_{2}$ ) occurred, the content of the bottom node remains unchanged.

## Input Control:

$\mathrm{I}_{1}$ : When $\rfloor(\uparrow \mathrm{L})$ is presented, the function block is executed.
$\mathrm{I}_{2}$ : error in.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=1$ (If the value of either top node or middle node is greater than 9999 , or $\mathrm{I}_{2}=1$ )
$\mathrm{O}_{3}$ : Overflow
$=1$, Product $\geqq 10000$
$=0$, Product $<10000$

## 【EXAMPLE】



## 【DESCRIPTION】

Let register（40005）＝25 and（40006）＝100．When contact 10007 is＇ ON ＇，the operation： $(40036)=(40005) \times(40006)$ is performed．

| 40005 | 0025 | $25 \times 100=2500$ |
| :---: | :---: | :---: |
| 40006 | 0100 | The product is stored in register 40007. |
| 40007 | 2500 |  |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  | $\bigcirc$ | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE |  |  | $\bigcirc$ | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1)0~65535
word $\times$ word $\rightarrow$ word (Binary)

## Description:

The value in the top node is multiplied by the value in the middle node, and the product is stored in the bottom node(long word).

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function output $\left(\mathrm{O}_{3}\right)$ may be used to determine whether or not an overflow has occurred .

## Node Description:

TOP: Multiplicand.
MIDDLE: Multiplier.
BOTTOM: Product.
Input Control:
$\mathrm{I}_{1}$ : When $\checkmark$
$(\uparrow L)$ is presented, the function block is executed.

## Function Output:

$$
\begin{aligned}
& \mathrm{O}_{1}=\mathrm{I}_{1} \\
& \mathrm{O}_{2}=0 \\
& \mathrm{O}_{3}: \text { Overflow } \\
& \quad=1, \text { Product } \geqq 65536 \\
& \quad=0, \text { Product }<65536
\end{aligned}
$$

## 【EXAMPLE】



## 【DESCRIPTION】

Let register（40005）＝9999 and（40006）＝2．When contact 10007 is＇ ON ＇，the operation： $(40036)=(40005) \times(40006)$ is performed．


SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE |  |  |  | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1)0~65535

Lword $\times$ Lword $\rightarrow$ Lword (Decimal)

## Description:

The value in the top node is multiplied by the value in the middle node, and the product is stored in the bottom node. All operands are long words.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function output $\left(\mathrm{O}_{3}\right)$ may be used to determine whether or not an overflow has occurred .

## Node Description:

TOP: Multiplicand, must be <= 99999999.
MIDDLE: Multiplier, must be <= 99999999.
BOTTOM :1. Product.
2.If error (refer to $\mathrm{O}_{2}$ ) occurred, the content of the bottom node remains unchanged.

## Input Control:

$\mathrm{I}_{1}$ : When $\Omega(\uparrow \square)$ is presented, the function block is executed.
$I_{2}$ : error in

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=1$ (If the value of either top node or middle node is greater than 99999999 , or $\mathrm{I}_{2}=1$ ).
$\mathrm{O}_{3}$ : Overflow
$=1$, Product $\geqq 100000000$
$=0$, Product $<100000000$

## 【EXAMPLE】



## 【DESCRIPTION】

Let register（40010／40011）＝12345 and $(40020 / 40021)=11$ ．When contact 10007 is＇ON＇，the operation：long word（40030）＝long word（40010）$\times$ long $\operatorname{word}(40020)$ is performed．

| 40010 | 0001 |
| :--- | :--- |
| 40011 | 2345 |
|  |  |
| 40020 | 0000 |
| 40021 | 0011 |
|  |  |
| 40030 | 0013 |
| 40031 | 5795 |
|  |  |
|  |  |

DECIMAL

## EIGHT DIGIT HEXADECIMAL MULTIPLIER

SYMBIL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE |  |  |  | $\bigcirc$ | (1) | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1)0~65535

Lword $\times$ Lword $\rightarrow$ Lword (Binary)

## Description:

The value in the top node is multiplied by the value in the middle node, and the product is stored in the bottom node. All operands are long words.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function output $\left(\mathrm{O}_{3}\right)$ may be used to determine whether or not an overflow has occurred .

## Node Description:

TOP: Multiplicand.
MIDDLE: Multiplier.
BOTTOM: Product.

## Input Control:

$\mathrm{I}_{1}$ : When $\Omega$
$(\uparrow L)$ is presented, the function block is executed.

## Function Output:

$$
\begin{aligned}
\mathrm{O}_{1} & =\mathrm{I}_{1} \\
\mathrm{O}_{2} & =0 \\
\mathrm{O}_{3} & : \text { Overflow } \\
& =1, \text { Product }>=4294967296\left(=100000000_{\text {hex }}\right) \\
& =0, \text { Product }<4294967296\left(=100000000_{\text {hex }}\right)
\end{aligned}
$$

## 【EXAMPLE】



## 【DESCRIPTION】

Let register（40010／40011）＝65535 and $(40020 / 40021)=11$ ．When contact 10007 is＇ON＇，the operation：long $\operatorname{word}(40030)=$ long $\operatorname{word}(40010) \times$ long $\operatorname{word}(40020)$ is performed．

| 40010 | 0000 | 40010 <br> 40011 | 0000 |
| :---: | :---: | :---: | :---: |
| 40011 | 65535 |  | FFFF |
| $\begin{aligned} & 40020 \\ & 40021 \end{aligned}$ | 0000 | $\begin{aligned} & 40020 \\ & 40021 \end{aligned}$ | 0000 |
|  | 0110 |  | 006E |
| 40030 | 0109 | $\begin{aligned} & 40030 \\ & 40031 \end{aligned}$ | 006D |
| 40031 | 65426 |  | FF92 |
|  | ECIMAL |  | ADEC |


| FMUL | FLOATING POINT MULTIPLIER | $\lrcorner \square$ | $\uparrow L$ |
| :---: | :---: | :---: | :---: |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE |  |  |  | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1)0~65535

## float $\times$ float $\rightarrow$ float

## Description:

The value stored in the top node is multiplied by the value in the middle node, and the product is stored in the bottom node. All operands are long words.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.

## Node Description:

TOP: Multiplicand.
MIDDLE: Multiplier.
BOTTOM: Product.

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark(\uparrow \square)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$
$\mathrm{O}_{3}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

Let register $(40010 / 40011)=5000$ and $(40020 / 40021)=2$ ．When contact 10025 is＇ ON ＇，the operation：long word（40030）＝long word（40010）$\times$ long word（40020）is performed．Function Output： $\mathrm{O}_{1}=\mathrm{ON}, \mathrm{O}_{2}=\mathrm{O}_{3}=\mathrm{OFF}$ ．


|  |  |  |  |
| :--- | :--- | :--- | :--- |
| DIV | FOUR DIGIT DECIMAL DIVIDER(1) | $\checkmark \square$ | $\uparrow \square$ |

## SYMBOL:



OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  | $\bigcirc$ | $\bigcirc$ | $\mathbb{1}$ | $\bigcirc$ |  |
| MIDDLE |  |  | $\bigcirc$ | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  | | (1)0~9999 |
| :--- | :--- |

Lword $\div$ word $\rightarrow$ word (Decimal)

## Description:

The value stored in the top node is divided by the value in the middle node, and the result is stored in the bottom node.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed, divisor is zero and overflow.

## Node Description:

TOP: Constant dividend, must be $<=9999$; else the LONG WORD value is used.
MIDDLE: Divisor, must be $<=9999$
BOTTOM: 1.Result of Division. The quotient is stored in the first word. Depending on the input control, the remainder or the first four digits after decimal point of quotient are stored in the second word.
2.If error occurred, the content of the bottom node remains unchanged.

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark(\uparrow \square)$ is presented, the function block is executed.
$\mathrm{I}_{2}=0$, the second word of the bottom node is used to store the remainder.
$=1$, the second word of the bottom node is used to store the first four digits after the decimal point. $\mathrm{I}_{3}$ : error in

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=1$, if overflow, i.e. quotient $>9999$
$\mathrm{O}_{3}($ error output $)=1$ (1.If the value of either top node or middle node is greater than 9999 or 2.If divisor $=0$ )

## 【EXAMPLE】



## (DESCRIPTION )

Let long word $(40090)=9999$ and $(40130)=10$. When contact 10005 is energized, $\mathrm{I}_{1}$ and $\mathrm{I}_{2}=' \mathrm{ON}$ '. The quotient $(=999)$ is stored in register 40053. Since $\mathrm{I}_{2}=$ ' ON ', thus the first four digits $(=9000)$ are stored in register 40054.

| 40053 | 0999 | Integer portion of the quotient <br> First four digits of the fractional portion of the quotient |  |
| :---: | :---: | :---: | :---: |
| 40054 | 9000 |  |  |
| 40090 | 0000 | Dividend | $9999 \div 10=999.9000$ |
| 40091 | 9999 |  |  |
| 40130 | 0010 | Divisor |  |

## DIVB $\quad$ EIGHT DIGIT HEXADECIMAL DIVIDER(1)

|  |  | DIVB |
| :---: | :---: | :---: |
|  | $\lrcorner \square$ | $\uparrow L$ |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  | $\bigcirc$ | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| MIDDLE |  |  | $\bigcirc$ | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1)0~65535

Lword $\div$ word $\rightarrow$ word (Binary)

## Description:

The value stored in the top node is divided by the value in the middle node, and the result is stored in the bottom node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed, divisor is zero and overflow.

## Node Description:

TOP: Dividend, (long word)
MIDDLE: Divisor
BOTTOM: 1.Result of Division. The quotient is stored in the first word. Depending on the input control, the remainder or the first four digits after decimal point of quotient are stored in the second word.
2.If error occurred, the content of the bottom node remains unchanged.

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark$
$(\uparrow L)$ is presented, the function block is executed.
$\mathrm{I}_{2}=0$, the second word of the bottom node is used to store the remainder.
$=1$, the second word of the bottom node is used to store the first four digits after the decimal point.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=1$, if overflow, i.e. quotient $>65535$
$\mathrm{O}_{3}($ error output $)=1$, if divisor $=0$.

## 【EXAMPLE】



## 【DESCRIPTION】

Let long word $(40090)=65535$ and $(40130)=12$ ．When contact 10005 is energized， $\mathrm{I}_{1}=' \mathrm{ON}$＇， and the quotient $(=5461)$ is stored in register 40053．Since $\mathrm{I}_{2}=$＇OFF＇，the remainder $(=0003)$ is stored in register 40054.


SYMBOL:


OPERANDS:

(1)0~9999
word $\div$ word $\rightarrow$ word (Decimal)

## Description:

The value stored in the top node is divided by the value in the middle node, and the result is stored in the bottom node.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed, divisor is zero and overflow.

## Node Description:

TOP: Dividend, must be <= 9999 .
MIDDLE: Divisor, must be <= 9999.
BOTTOM :1.Result of Division. The quotient is stored in the first word. Depending on the input control, the remainder or the first four digits after decimal point of quotient are stored in the second word.
2.If error occurred, the content of the bottom node remains unchanged.

## Input Control:

$\mathrm{I}_{1}$ : When $\square(\uparrow L)$ is presented, the function block is executed.
$I_{2}=0$, the second word of the bottom node is used to store the remainder.
$=1$, the second word of the bottom node is used to store the first four digits after the decimal point.
$\mathrm{I}_{3}=$ error in

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$
$\mathrm{O}_{3}$ (error output)=1(1. If divisor $=0$ or
2. If the value of either top node or middle node is greater than 9999. )

## 【EXAMPLE】



## 【DESCRIPTION】

Let long $\operatorname{word}(40090)=9999$ and $(40130)=10$ ．When contact 10005 is energized，$I_{1}$ and $\mathrm{I}_{2}=$＇ON＇．The quotient $(=999)$ is stored in register 40053．Since $\mathrm{I}_{2}=$＇ON＇，thus the first four digits（＝9000）are stored in register 40054.

$$
9999 \div 10=999.9000
$$

| 40053 | 0999 | Quotient |
| :---: | :---: | :---: |
| 40054 | 9000 | First four digits of the fractional portion of the quotient |
| 40090 | 9999 | Dividend |
| 40130 | 0010 | Divisor |


| DVBM | FOUR DIGIT HEXADECIMAL DIVIDER | $\checkmark L$ | $\uparrow\llcorner$ |
| :---: | :---: | :--- | :--- |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  | $\bigcirc$ | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| MIDDLE |  |  | $\bigcirc$ | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1)0~65535
word $\div$ word $\rightarrow$ word (Binary)

## Description:

The value stored in the top node is divided by the value in the middle node, and the result is stored in the bottom node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed, divisor is zero and overflow.

## Node Description:

TOP: Dividend.
MIDDLE: Divisor.
BOTTOM: 1.Result of Division. The quotient is stored in the first word. Depending on the input control, the remainder or the first four digits after decimal point of quotient are stored in the second word.
2.If error occurred, the content of the bottom node remain unchanged.

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark \square$ ( $\uparrow$ ) is presented, the function block is executed.
$\mathrm{I}_{2}=0$, the second word of the bottom node is used to store the remainder.
$=1$, the second word of the bottom node is used to store the first four digits after the decimal point.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$
$\mathrm{O}_{3}($ error output $)=1$, if divisor $=0$.

## 【EXAMPLE】



## 【DESCRIPTION】

Let long word $(40090)=65535$ and $(40130)=12$ ．When contact 10005 is energized， $\mathrm{I}_{1}=$＇ON＇， and the quotient $(=5461)$ is stored in register 40053．Since $\mathrm{I}_{2}=$＇OFF＇，the remainder $(=0003)$ is stored in register 40054.

| 40053 | 5461 | Quotient |
| :--- | :--- | :--- |
| 40054 | 0003 | Remainder |$\quad 65535 \div 12=5461$, remainder 3

OPERANDS:


## Description:

The value stored in the top node is divided by the value in the middle node, and the result is stored in the bottom node. Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not. Function outputs can be used to determine whether the function block has been executed, divisor is zero and overflow.

## Node Description:

TOP: Dividend, must be <= 99999999.
MIDDLE: Divisor, must be <= 99999999.
BOTTOM:1.Result of Division. The quotient is stored in the first and second words. Depending on the input control, the remainder or the first eight digits after decimal point of quotient are stored in the third and fourth words.
2.If error occurred, the content of the bottom node remains unchanged.

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark(\neg \square)$ is presented, the function block is executed.
$\mathrm{I}_{2}=0$, the third and the fourth words of the bottom node is used to store the remainder.
$=1$, the third and the fourth words of the bottom node is used to store the first eight digits after the decimal point.
$\mathrm{I}_{3}$ : error in

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$
$\mathrm{O}_{3}$ (error output) $=1$ (1. If divisor $=0$ or
$\quad \begin{aligned} & \text { 2. If the value of either top node or middle node is greater than 99999999) }\end{aligned}$

## 【EXAMPLE】



## 【DESCRIPTION】

Let long word $(40090)=99999999$ and long $\operatorname{word}(40130)=11$ ．When contact 10005 is energized， $\mathrm{I}_{1}=$＇ ON ＇，and the quotient $(=9090910)$ is stored in the long word 40053．Since $\mathrm{I}_{2}=$ ＇OFF＇，the remainder $(=0001)$ is stored in register 40055 and 40056.

| 40053 | 0909 | $\longrightarrow \text { Quotient(=9090910) }$ |
| :---: | :---: | :---: |
| 40054 | 0910 |  |
| 40055 | 0000 | Remainder $(=0000)$ |
| 40056 | 0001 |  |
| 40090 | 9999 | $\longrightarrow$ Dividend（＝99999999） |
| 40091 | 9999 |  |
| 40130 | 0000 | $\longrightarrow$ Divisor（＝11） |
| 40131 | 0011 |  |
|  |  |  |

OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  |  | $\bigcirc$ | (1) | $\bigcirc$ |  |
| MIDDLE |  |  |  | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1)0~65535

Lword $\div$ Lword $\rightarrow$ Lword

## Description:

The value stored in the top node is divided by the value in the middle node, and the result is stored in the bottom node.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed, divisor is zero and overflow.

## Node Description:

TOP: Dividend.
MIDDLE: Divisor.
BOTTOM: Result of Division. The quotient is stored in the first and second words. Depending on the input control, the remainder or the eight digits after decimal point of quotient are stored in the third and fourth words.

## Input Control:

$\mathrm{I}_{1}$ : When $\sqrt{ }$ ( $\uparrow \quad$ ) is presented, the function block is executed.
$\mathrm{I}_{2}=0$, the third and the fourth words of the bottom node is used to store the remainder.
$=1$, the third and the fourth words of the bottom node is used to store the first eight digits after the decimal point.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$
$\mathrm{O}_{3}=1$, if divisor $=0$.

## 【EXAMPLE】



## 【DESCRIPTION】

Let long $\operatorname{word}(40090)=305419896$ and long word $(40130)=1100$ ．When contact 10005 is energized， $\mathrm{I}_{1}=$＇ ON ＇，and the quotient（ $=277654$ ）is stored in long word 40053．Since $\mathrm{I}_{2}=$＇OFF＇， the remainder $(=0000)$ is stored in register 40055 and 40056.

| 40053 | 0004 | $\rightarrow$ Quotient $=277654$ ） |
| :---: | :---: | :---: |
| 40054 | 3C96 |  |
| 40055 | 0000 | Remainder $(=0000)$ |
| 40056 | 0000 |  |
| 40090 | 1234 | $\rightarrow$ Dividend（ $=305419896$ ） |
| 40091 | 5678 |  |
| 40130 | 0000 | $\rightarrow$ divisor $=1100$ ） |
| 40131 | 044C |  |


| FDIV | FLOATING POING DIVIDER | $\checkmark\llcorner$ | $\uparrow\llcorner$ |
| :---: | :---: | :---: | :---: |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  |  | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| MIDDLE |  |  |  | $\bigcirc$ | $(1)$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |

(1)0~65535

## float $\div$ float $\rightarrow$ float

## Description:

The value stored in the top node is divided by the value in the middle node, and the result is stored in the bottom node.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed, divisor is zero and overflow.

## Node Description:

TOP: Dividend.
MIDDLE: Divisor.
BOTTOM: Quotient.

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark L(\uparrow L)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$
$\mathrm{O}_{3}=1$, if divisor $=0$.

## 【EXAMPLE】



## 【DESCRIPTION】

Let register（40010／40011）＝5000 and（40020／40021）＝2．When contact 10025 is＇ON＇，the operation：long word（40030）＝long word（40010）$\div$ long $\operatorname{word}(40020)$ is performed．Function Output： $\mathrm{O}_{1}=\mathrm{ON}, \mathrm{O}_{2}=\mathrm{O}_{3}=\mathrm{OFF}$ ．


## ISQR

| ISQR | SQUARE ROOT OF AN INTEGER | $\lrcorner \square$ | $\uparrow\llcorner$ |
| :---: | :---: | :--- | :--- |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  | $O$ | $\bigcirc$ |  | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
|  |  |  |  |  |  |  |  |

## Description:

The square root of the value stored in the top node is found and stored in the bottom node. The result of the square root operation is truncated to integer. Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not. Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: An integer whose square root is desired.
BOTTOM: Square root.

## Input Control:

$\mathrm{I}_{1}$ : When $\left.\rfloor( \lrcorner \square\right)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

Let $(40120)=400$ ．When this rung is scanned，the square root of the values stored in the top node are stored in the bottom node．


| FSQR | SQUARE ROOT OF A FLOATING POINT <br> NUMBER | $\lrcorner L$ | $\uparrow\llcorner$ |
| :---: | :---: | :--- | :--- |

SYMBOL:


OPERANDS:


## Description:

The square root of the value stored in the top node is found and stored in the bottom node. Input control $\left(I_{1}\right)$ is used to determine whether this function block is to be executed or not. Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: A floating point number whose square root is desired.
BOTTOM: Square root.

## Input Control:

$\mathrm{I}_{1}$ : When $\square$ ( $\uparrow$ ) is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

Let $(40120 / 40121)=2500$ ．When this rung is scanned，the square root of the values stored in the top node are stored in the bottom node（40130／40131）．



## Description:

The content of the top node is filled onto the table defined in the following address(es) by the middle node.
Table length is defined in the bottom node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Input control ( $\mathrm{I}_{2}$ ) is used to define the action mode of the INDEX.
Input control ( $\mathrm{I}_{3}$ ) is used to clear the INDEX.
Function outputs can be used to determine whether the function block has been executed and whether the INDEX exceeded the table length.

## Node Description:

TOP: Source register.
MIDDLE: Reference register. First word defined as INDEX into the target table. If the value of the INDEX is equal to zero, then the INDEX is pointing to the first entry in the target table. The target table starts with the second word.
BOTTOM: Table Length. If the INDEX value is greater than or equal to this number, table movement is prohibited disregarding the state of $\mathrm{I}_{1}$.
Input Control:
$\mathrm{I}_{1}$ : When $\rfloor(\uparrow L)$ is presented, the function block is executed.
$\mathrm{I}_{2}$ : INDEX control.
$=0$, INDEX is incremented by one after each execution.
$=1$, INDEX remains unchanged.
$\mathrm{I}_{3}$ : Reset INDEX.
$=1$, clear INDEX to 0 .

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}$ : INDEX indicator.
$=1$, INDEX $\geqq$ table length, the INDEX is pointing to an address beyond table limit.
$\mathrm{O}_{3}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10005 is energized，the content of input register 40120 is copied to table registers （40101～40110），one register per scan．During the action INDEX（40100）increments by one after each scan．

When INDEX（in 40100）reaches preset value of BOTTOM node（\＃00010），then coil 00001 is energized and the content of register 40100 is cleared．The movement continues until contact 10005 is OFF．


BEFORE

## AFTER



## Description:

The content of the top node is moved to the following address(es) defined by the middle node. Table length is defined in the bottom node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Input control ( $\mathrm{I}_{2}$ ) is used to define the action mode of the INDEX.
Input control ( $\mathrm{I}_{3}$ ) is used to clear the INDEX.
Function outputs can be used to determine whether the function block has been executed and whether the INDEX exceeded the table length.

## Node Description:

TOP: Source table.
MIDDLE: Source INDEX is defined at the first word. If the value of the INDEX is equal to zero, then the INDEX is pointing to the first entry in the source table. The target register is in the second word.
BOTTOM: Table Length. If the INDEX value is greater than or equal to this number, table movement is prohibited disregarding the state of $\mathrm{I}_{1}$.

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark(\uparrow L)$ is presented, the function block is executed.
$\mathrm{I}_{2}$ : INDEX control.
$=0$, INDEX is incremented by one after each execution.
$=1$, INDEX remains unchanged.
$\mathrm{I}_{3}$ : Reset INDEX.
$=1$, clear INDEX to 0 .
Function Output:
$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}$ : INDEX indicator.
$=1$, INDEX $\geqq$ table length, the INDEX is pointing to an address beyond table limit.
$\mathrm{O}_{3}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10054 is energized，source data pointed to by INDEX（40101）is moved to 40201 （the next address defined by the middle node）．For every scan of the PLC controller，data movement occurs once until the INDEX reaches the end of table（\＃00050）．Then Coil 00129 is energized and the INDEX is cleared．In this manner，data movement can be repeated．The following is the state after the nth scan since INDEX reset to 0 ．


Table

| T-> T | MOVE FROM ONE TABLE TO ANOTHER TABLE | $\checkmark$ ¢ | $\dagger$ ¢ |
| :---: | :---: | :---: | :---: |

SYMBOL:


OPERANDS:

(1)1~255

## Description:

The content of the top node is moved to the following address(es) defined by the middle node. Table length is defined in the bottom node. Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not. Input control ( $\mathrm{I}_{2}$ ) is used to define the action mode of the INDEX. Input control $\left(\mathrm{I}_{3}\right)$ is used to clear the INDEX. Function outputs can be used to determine whether the function block has been executed and whether the INDEX exceeded the table length.

## Node Description:

TOP: Source table.
MIDDLE: INDEX is defined at the first word. If the value of the INDEX is equal to zero, then the INDEX is pointing to the first entry in the target register. The target table starts at the second word. BOTTOM: Table Length. If the INDEX value is greater than or equal to this number, table movement is prohibited disregarding the state of $\mathrm{I}_{1}$.

## Input Control:

$\mathrm{I}_{1}$ : When $\square(\uparrow \square)$ is presented, the function block is executed.
$\mathrm{I}_{2}$ : INDEX control.
$=0$, INDEX is incremented by one after each execution.
$=1$, INDEX remains unchanged.
$\mathrm{I}_{3}$ : Reset INDEX.
$=1$, clear INDEX to 0 .
Function Output:
$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}$ : INDEX indicator.
$=1$, INDEX $\geqq$ table length, the INDEX is pointing to an address beyond table limit.
$\mathrm{O}_{3}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 is energized and after six scan cycles，data stored in registers 30001～30006 are moved to registers 40101～40106．When the INDEX in the middle node $=6$ ，the coil 00001 is energized and the INDEX is cleared．Data movement continues until contact 10001 is OFF．


After 6 scans


## Description:

This function compares the tables pointed by the top and middle node. If a difference is found between the corresponding table locations, then that the index of that element is stored in the middle node. Table length is in the bottom node.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Input control ( $\mathrm{I}_{2}$ ) is used to define the action mode of the INDEX.
Function outputs can be used to determine whether the function block has been executed and whether those tables are different or not.

## Node Description:

TOP: Starting address of the first table.
MIDDLE: INDEX is defined in the first word. The starting address of the second table is defined in the second word. The INDEX has two uses. First, it is used to indicate the index of the difference data after the last comparison. Second, it is used to indicate the next comparison will start below this index.
BOTTOM: Table length.

## Input Control:

$\mathrm{I}_{1}$ : When $\square(\uparrow L)$ is presented, the function block is executed. Based on the state of $\mathrm{I}_{2}$, comparison continues until a difference is found or the end of table is reached. When execution terminates, the INDEX points to the address where the difference is found or zero (if no difference).
$\mathrm{I}_{2}$ : Action mode of INDEX.
$=0$, start comparison from the next address pointed to by the INDEX.
$=1$, start comparison from the first entry of the table.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=1$, if tables are different.
$\mathrm{O}_{3}$ is the complement of $\mathrm{O}_{2}$.

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10005 is energized，the table starting from 40211 is compared against the table starting from 40451 ．Since the $6^{\text {th }}$ entries in both tables are different，then register $40450=00006$ ， and coil 00001＝＇ON＇．



## Description:

Reference value is defined in the second word of the middle node. Search starts from the table defined in the top node. If the same value as the reference is found, then this function stops and the INDEX where the same value was found is stored in the middle node. Table length is defined in the bottom node. Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Input control ( $\mathrm{I}_{2}$ ) is used to define the action mode of the INDEX.
Function outputs can be used to determine whether the function block has been executed and whether a same value with the reference has been found.

## Node Description:

TOP: Starting address of the table.
MIDDLE: INDEX is defined in the first word. The reference value is defined in the second word of the middle node. The INDEX has two uses. First, it is used to indicate the index of the same data after the last search. Second, it is used to indicate the next search will start below this index.
BOTTOM: Table length.

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark(\uparrow L)$ is presented, the function block is executed. Based on the state of $\mathrm{I}_{2}$, search continues until a match is found or the end of the table is reached. When execution terminates, the INDEX points to the address where the difference is found.
$\mathrm{I}_{2}$ : Action mode of INDEX.
$=0$, start search from the next address pointed to by the INDEX.
$=1$, start search from the first entry of the table.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=1$, if tables are different.

## 【EXAMPLE】



## 【DESCRIPTION】

When $\mathrm{I} 1=1$ ，search starts from the table starting from 40211．INDEX（40206）reset to 0 at the beginning．If a table entry is found to be the same as the value of 40207 （next address defined in the middle node），in this example，（40216）＝00666＝（40207），then 00006 is stored in the middle node，and coil $00010=$＇ON＇．



## Description:

Using register as a unit, this function performs table rotate or shift. The table is defined in the top node.
Table length is defined in the bottom node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Input control ( $\mathrm{I}_{2}$ ) is used to define the direction.
Input control ( $\mathrm{I}_{3}$ ) is used to define the mode.
Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: Starting address of the table.
MIDDLE: Table after processing.
BOTTOM :Table length.
LEFT


## SHIFT MODE

## Input Control:

$\mathrm{I}_{1}$ : Execution control.
When $\checkmark \square(\uparrow L)$ is presented, the function block is executed.
$\mathrm{I}_{2}$ : Direction.
$=0$, Left.
$=1$, Right.
$\mathrm{I}_{3}$ : mode.

$$
\begin{aligned}
& =0 \text {, Shift. } \\
& =1 \text {, Rotate. }
\end{aligned}
$$

## Function Output:

$$
\begin{aligned}
& \mathrm{O}_{1}=\mathrm{I}_{1} \\
& \mathrm{O}_{2}=0 \\
& \mathrm{O}_{3}=0
\end{aligned}
$$



ROTATE MODE

LEFT


## 【EXAMPLE 1】



## 【DESCRIPTION】

When contact 10001 receives a transition from＇ OFF ＇to＇ ON ＇and $\mathrm{I}_{2}=\mathrm{I}_{3}=1$ ，then a right rotate operation is performed．


【EXAMPLE 2】


## 【DESCRIPTION】

BEFORE

|  |  |
| :--- | :--- |
| 40010 | 1111 |
| 40011 | 2222 |
| 40012 | 3333 |
|  |  |


| AFTER |  |
| :---: | :---: |
| 40010 | 2222 |
| 40011 | 3333 |
| 40012 | 1111 |
|  |  |


| TXHG | TABLE EXCHANGE | $\rfloor L$ | $\uparrow\llcorner$ |
| :---: | :---: | :---: | :---: |

## SYMBOL:



OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| TOP | O |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE | O |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| BOTTOM |  |  |  |  | $(1)$ |  |  |

(1)1~255

## Description:

Entries exchange between two tables. Tables are defined in the top and middle nodes. Table length is defined in the bottom node.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: The first table.
MIDDLE: The second table.
BOTTOM: Table length.
Input Control:
$\mathrm{I}_{1}$ : When $\curvearrowleft(\uparrow \square)$ is presented, the function block is executed.

## Function Output:

$$
\begin{aligned}
& \mathrm{O}_{1}=\mathrm{I}_{1} \\
& \mathrm{O}_{2}=0 \\
& \mathrm{O}_{3}=0
\end{aligned}
$$

## 【EXAMPLE】

| 10001 | 40010 <br> 40020 <br> TXHG <br> $\# 00005$ |
| :---: | :---: |

## 【DESCRIPTION】

When 10001 receives a transition from＇OFF＇to＇ON＇，then the entries of the two tables as defined in the top and middle nodes are swapped．

| 40010 | 01111 | 40020 | 06666 |
| :---: | :---: | :---: | :---: |
| 40011 | 02222 | 40021 | 07777 |
| 40012 | 03333 | 40022 | 08888 |
| 40013 | 04444 | 40023 | 09999 |
| 40014 | 05555 | 40024 | 00000 |
|  |  |  |  |

BEFORE


AFTER

## BLKM

| BLKM | MEMORY BLOCK MOVE | $\lrcorner L$ | $\uparrow L$ |
| :---: | :---: | :--- | :--- |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| TOP | O | O | O | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE | O |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| BOTTOM |  |  |  |  | $(1)$ |  |  |

$$
\text { (1) } 1 \sim 255
$$

## Description:

Memory contents of the table defined in the top node are copied to the table defined in the middle node in one scan. Table length is defined in the bottom node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: Starting address of the source memory block.
MIDDLE: Starting address of the target memory block.
BOTTOM: Block length.
Input Control:
$\mathrm{I}_{1}$ : When $\left.\rfloor( \lrcorner \square\right)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$
$\mathrm{O}_{3}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

When 10001 receives a transition from＇OFF＇to＇ON＇，then the entries of the first tables as defined in the top node（40211）are moved to the second table defined in the middle node（40451）．

| 40211 | 01111 |
| :---: | :---: |
| 40212 | 02222 |
| 40213 | 03333 |
| 40214 | 04444 |
| 40215 | 05555 |
| 40216 | 06666 |
| 40217 | 07777 |
| 40218 | 08888 |
| 40219 | 09999 |
| 40220 | 00000 |
|  |  |


|  |  |
| :--- | :--- |
|  | 40451 |
| 40452 | 00000 |
|  | 00000 |
| 40453 | 00000 |
|  | 00454 |
| 40455 | 00000 |
|  | 00000 |
| 40456 | 00000 |
|  | 000000 |
| 40458 | 00000 |
| 40459 | 00000 |
| 40460 | 00000 |
| 40461 | 00000 |

BEFORE


## AFTER

|  | PUSH |  |  |
| :---: | :---: | :---: | :---: |
| PUSH | PUSH FROM REGISTER TO STACK | $\boxed{\square}$ | $\uparrow$ |

## SYMBOL:



OPERANDS:

(1) 1~255

## Description:

This function pushes the content of a register (TOP NODE) to a predefined stack(MIDDLE NODE).
Stack size is defined in the BOTTOM NODE.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Input control ( $\mathrm{I}_{2}$ ) is used to define the stack mode.
Function outputs can be used to determine whether the function block has been executed, and whether the stack is full.

## Node Description:

TOP: Data to be pushed into stack.
MIDDLE: INDEX is defined in the first word. The starting of the stack is defined in the second word of the middle node. If the INDEX is equal to zero, then the INDEX is pointing to the top of the stack.

BOTTOM: Length of stack.

## Input Control:

$\mathrm{I}_{1}$ : When $\square(\uparrow \quad)$ is presented, the function block is executed. INDEX is incremented by 1.
$\mathrm{I}_{2}=0$, data is pushed into the stack at designated address. (LIFO). Used as STACK.
$=1$, data is put into the bottom of the stack, the original content at the top of the stack is moved to the next word. (FIFO). Used as QUEUE.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}:=1$, if stack is full, i.e. $\mathrm{INDEX}=$ stack length.
$\mathrm{O}_{3}=0$

## 〔EXAMPLE 1】

| H | $\uparrow 40601$ |
| :---: | :---: |
| 00015 |  |
|  | 40500 |
|  | PUSH |
|  | \＃00005 |

## 【DESCRIPTION】

Since $\mathrm{I}_{2}=$＇OFF＇，thus the operation mode is LIFO（Last In First Out）．
When $\mathrm{I}_{1}$ receives an＂OFF $\rightarrow \mathrm{ON}$＂transition，the PUSH function is executed as follows：


For the next＂OFF $\rightarrow \mathrm{ON}$＂transition on $\mathrm{I}_{1}$ ：


## 【EXAMPLE 2】

| 00015 | 40601 <br> 40500 <br> PUSH <br> $\#$ <br> $\# 00005$ |
| ---: | :--- |

【DESCRIPTION】Since $\mathrm{I}_{2}=$＇ON＇，thus the operation mode is FIFO（First in first out）． When $\mathrm{I}_{1}$ receives an＂OFF $\rightarrow \mathrm{ON}$＂transition，the PUSH function is executed as follows：


BEFORE
AFTER

For the next＂OFF $\rightarrow \mathrm{ON}$＂transition on $\mathrm{I}_{1}$ ：


BEFORE
AFTER

| POP | POP FROM STACK TO REGISTER | $\rfloor \square$ | $\uparrow\llcorner$ |
| :---: | :---: | :--- | :--- |

SYMBOL:


OPERANDS:

(1) 1~255

## Description:

This function moves the content of a stack defined in the top node to the register defined in the middle node.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed, and whether the stack is empty.

## Node Description:

TOP: INDEX is defined in the first word. The starting of the stack is defined in the second word of the middle node. If the INDEX is equal to zero, then the INDEX is pointing to the top of the stack.
MIDDLE: Data retrieved from stack.
BOTTOM: stack length.

## Input Control:

$\mathrm{I}_{1}$ : When $\lceil(\uparrow \square)$ is presented, the function block is executed. INDEX is decremented by 1 first, then the data is retrieved according to the INDEX.

## Function Output:

```
O
O
O
    =1, if stack is empty, i.e. INDEX =0.
```


## 【EXAMPLE】

| 10020 | 40340 <br> 40401 <br> POP <br> $\# 00100$ |
| :---: | :--- |

## 【DESCRIPTION】

When contact $10020=$＇ON＇，the INDEX（40340）is decremented by 1 ；then the value pointed by the INDEX is retrieved and stored in the location pointed by the middle node．Through repeated conducting of contact 10020，the values in the stack are POPed successively．


BEFORE
AFTER

| AND | AND OPERATION FOR ARRAYS | $\downharpoonleft\llcorner$ | $\uparrow\llcorner$ |
| :---: | :---: | :--- | :---: |

## SYMBOL:



OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP | $\bigcirc$ | $\bigcirc$ | $\bigcirc$ | $\bigcirc$ | $(1$ | $\bigcirc$ |  |
| MIDDLE | $\bigcirc$ |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| BOTTOM |  |  |  |  | $\square$ |  |  |

## (1) 0~65536

—1~255

## Description:

The contents of top and middle nodes are ANDed, and the result is stored in the middle node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed.
Remark: When the content of the top node is a constant, the constant and the contents of middle node are ANDed and the result is stored in the middle node.

## Node Description:

TOP: Source Array 1, or constant.
MIDDLE: Source Array 2, Resultant Array (after processing).
BOTTOM: Length of Array.

## Input Control:

$\mathrm{I}_{1}$ : When $\sqrt[\square]{ }(\uparrow)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$
$\mathrm{O}_{3}=0$

## 【EXAMPLE 1】



## 【DESCRIPTION】

When contact 10025 is energized，the contents of registers 40090 and 40095 are ANDed，and the result is returned to register 40095.


## 【EXAMPLE 2】

－10025 | \＃0F0Fh |
| :---: |
| 40095 |
| AND |
| \＃00002 |

| 40095 | AAAA | AND \＃0F0FH | 0A0A |
| :---: | :---: | :---: | :---: |
| 40096 | 0000 | AND \＃0F0FH | 0D0D |
|  |  |  |  |


|  |  |  |  |
| :---: | :---: | :---: | :---: |
| OR OR |  |  |  |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP | $\bigcirc$ | $\bigcirc$ | $\bigcirc$ | $\bigcirc$ | $(1$ | $\bigcirc$ |  |
| MIDDLE | $\bigcirc$ |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| BOTTOM |  |  |  |  | $\square$ |  |  |

## (1) $0 \sim 65535$

—1~255

## Description:

The contents of top and middle nodes are ORed, and the result is stored in the middle node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed.
Remark: When the content of the top node is a constant, the constant and the contents of middle node are ORed and the result is stored in the middle node.

## Node Description:

TOP: Source Array 1, or constant.
MIDDLE: Source Array 2, Resultant Array (after processing).
BOTTOM: Length of Array.
Input Control:
$\mathrm{I}_{1}$ : When $\square(\uparrow L)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$
$\mathrm{O}_{3}=0$

## 【EXAMPLE 1】

10025 | 40090 |
| :---: |
| 40095 |
| OR |
| $\# 00002$ |

## 【DESCRIPTION】

When contact 10025 is energized，the contents of registers 40090 and 40095 are ORed，and the result is returned to register 40095.


## 【EXAMPLE 2】



40095

| BBBB | AND＇AAAA＇ | BBBB | 40095 |
| :---: | :---: | :---: | :---: |
| CCCC | $\xrightarrow{\text { AND＇AAAA＇}}$ | EEEE | 40096 |


| XOR | XOR OPERATION FOR ARRAYS | $\checkmark L$ | $\uparrow L$ |
| :---: | :---: | :--- | :--- |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP | O | O | O | $\bigcirc$ | (1) | ○ |  |
| MIDDLE | O |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| BOTTOM |  |  |  |  | $\square$ |  |  |

(1) 0~65535

1~255

## Description:

The contents of top and middle nodes are XORed, and the result is stored in the middle node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed.
Remark: When the content of the top node is a constant, the constant and the contents of middle node are XORed and the result is stored in the middle node.

## Node Description:

TOP: Source Array 1 or constant.
MIDDLE: Source Array 2, Resultant Array (after processing).
BOTTOM: Length of Array.
Input Control:
$\mathrm{I}_{1}$ : When $\left.\rfloor( \rfloor \square\right)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$
$\mathrm{O}_{3}=0$

## 【EXAMPLE 1】

|  | 40090 <br> 40095 <br> XOR <br> \＃00002 |
| ---: | :--- | :--- |
| 10025 |  |
|  |  |

## 【DESCRIPTION】

When contact 10025 is energized，the contents of registers 40090 and 40095 are XORed，and the result is returned to register 40095.


## 【EXAMPLE 2】

10025 | \＃5555h |
| :--- |
| 40095 |
| XOR |
| $\# 00002$ |

| 40095 | BBBB | XOR＇5555＇ | EEEE | 40095 |
| :---: | :---: | :---: | :---: | :---: |
| 40096 | CCCC | XOR＇5555＇ | 9999 | 40096 |
|  |  |  |  |  |

## COMP

1'S COMPLEMENT FOR ARRAYS

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP | O | O | O | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE | $\bigcirc$ |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| BOTTOM |  |  |  |  | © |  |  |

(1) 1~255

## Description:

1 's complement is obtained for the content of the top node, and the result is stored in the middle node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: Source Array.
MIDDLE: Resultant Array.
BOTTOM: Length of Array.
Input Control:
$\mathrm{I}_{1}$ : When $\square$ ) is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$
$\mathrm{O}_{3}=0$

## 【EXAMPLE】

| 10025 | 40090 <br> 40095 <br> COMP <br> $\# 00001$ |
| :---: | :---: |

## 【DESCRIPTION】

When contact 10025 is energized， 1 ＇s complement is obtained for the content of register 40090， and the result is returned to register 40095.


|  |  | CMPR |  |
| :---: | :--- | :---: | :---: |
| CMPR | BIT COMPARISON BETWEEN TWO MATRIX | $\boxed{\square}$ | $\uparrow \quad$ |

## SYMBOL:



OPERANDS:

(1) 1~255

## Description:

This function compares the matrix pointed by the top and middle nodes. If a difference is found between the corresponding matrix locations, then the index of that element is stored in the middle node. Input control $\left(\mathrm{I}_{1}\right)$ is used to determine whether this function block is to be executed or not. Input control $\left(\mathrm{I}_{2}\right)$ is used to indicate the position where comparison is started. Function outputs can be used to determine whether the function block has been executed and whether those tables are different or not.


## Node Description:

TOP: matrix 1.
MIDDLE: Index and matrix 2. INDEX is stored in the first word. Matrix 2 is stored starting from the second word. If the value of the INDEX is zero after searching, it represents that the contents of the two matrixes are identical.
BOTTOM: Length of matrix (word).

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark(\uparrow L)$ is presented, the function block is executed. When a difference is found, the INDEX points to the position where the difference is found.
$\mathrm{I}_{2}$ : Start position of the comparison.
$=0$, start from the position pointed to by the INDEX.
$=1$, start from the first position

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=1$, if a difference is found.
$\mathrm{O}_{3}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 is energized，the matrix starting from 40010 is compared against the matrix starting from 40100 ．Since the fourth bit is different，then the index of that location is stored in the middle node and coil 00002 is energized．



## Description:

Using the bit as a unit, this function performs array rotate or shift. The result is stored in the middle node. The matrix is defined in the top node.
Table length is defined in the bottom node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Input control ( $\mathrm{I}_{2}$ ) is used to define the direction.
Input control ( $\mathrm{I}_{3}$ ) is used to define the mode.
Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: Source matrix.

## SHIFT MODE

MIDDLE: Target matrix.
BOTTOM: Length of matrix (word).
LEFT


## Input Control:

$\mathrm{I}_{1}$ : Execution control.
When $\left.\quad L_{( } \uparrow\right)$ is presented,

## RIGHT


rotate/shift operation is performed one bit per scan.
$\mathrm{I}_{2}$ : Direction
$=0$, Left.
$=1$, Right.
$\mathrm{I}_{3}$ : Mode.
$=0$, Shift.
$=1$, Rotate.
Function Output:
$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}$ : Bit shifted into this position.
$\mathrm{O}_{3}=0$

RIGHT

## ROTATE MODE

LEFT


## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 receives a transition from＇OFF＇to＇ ON ＇and $\mathrm{I}_{2}=\mathrm{I}_{3}=1$ ，then a right rotate operation is performed．


| ODSR | NIBBLE ROTATE/SHIFT FOR MATRIX | $\checkmark\llcorner$ | $\uparrow \mathrm{L}$ |
| :---: | :--- | :--- | :--- |

## SYMBOL:



OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP | O | O | O | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE | $\bigcirc$ |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| BOTTOM |  |  |  |  | $(1)$ |  |  |

(1) 1~255

## Description:

Using the nibble as a unit, this function performs array rotate or shift. The result is stored in the middle node. The matrix is defined in the top node. Table length is defined in the bottom node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Input control ( $\mathrm{I}_{2}$ ) is used to define the direction.
Input control ( $\mathrm{I}_{3}$ ) is used to define the mode.
Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: Source matrix
MIDDLE: Target matrix
BOTTOM: Length of matrix (word)
Input Control:
$\mathrm{I}_{1}$ : Execution control.

When $\quad \square(\uparrow \square)$ is presented,
rotate/shift operation is performed one nibble per scan.
n: Direction

## ROTATE MODE

$=0$, Left
=1, Right
LEFT

$=0$, Shift.
$=1$, Rotate.
RIGHT


[^0]
## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 receives a transition from＇OFF＇to＇ ON ＇and $\mathrm{I}_{2}=\mathrm{I}_{3}=1$ ，then a right rotate operation is performed．

Initial State | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | 0

After 1st Scan | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

After 2nd Scan | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |



## Description:

This function is used to SET or CLEAR a certain bit in a matrix. Bit location is defined in the top node.
Array to be modified is defined in the middle node. Array length (WORD)is defined in the bottom node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Input control ( $\mathrm{I}_{2}$ ) is used to define the action (SET or Clear).
Input control $\left(\mathrm{I}_{3}\right)$ is used to define the behavior of the INDEX.
Function outputs can be used to determine whether the function block has been executed, and the status of the INDEX.

## Node Description:

TOP: INDEX (pointing to the bit to be modified). INDEX $=1 \rightarrow$ The first bit.
MIDDLE: Source matrix.
BOTTOM: Matrix length (word).

## Input Control:

$\mathrm{I}_{1}$ : When $\lceil(\uparrow \square)$ is presented, the function block is executed.
$\mathrm{I}_{2}$ : Action.
$=0$, bit clear
$=1$, bit set
$\mathrm{I}_{3}$ : INDEX control. If $\mathrm{I}_{3}=1$ and TOP $=4 \mathrm{xxxx}$, then the INDEX is incremented by 1 after execution.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=\mathrm{I}_{2}$
$\mathrm{O}_{3}$ : Status of the INDEX.
$=1$, if INDEX is larger than the value of the BOTTOM node times 16 .

## 【EXAMPLE】



## 【DESCRIPTION】

00005 is stored in the top node（40093）．When contact 10001 is energized，and $\mathrm{I}_{2}=1$ ，then the $5^{\text {th }}$ bit of the matrix starting from 40733 to 40735 is set to 1 ．


| SENS | SENSING OF A BIT IN MATRIX | $\rfloor\llcorner$ | $\uparrow \square$ |
| :---: | :---: | :---: | :---: |

## SYMBOL:



OPERANDS:

(1) 1~255
(2) 1~65535

## Description:

This function is used to sense a certain bit in a matrix. Bit location is defined in the top node. Matrix to be modified is defined in the middle node. Array length is defined in the bottom node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Input control ( $\mathrm{I}_{2}$ ) is used to define the behavior of the INDEX.
Input control ( $\mathrm{I}_{3}$ ) is used to reset the INDEX.
Function outputs can be used to determine whether the function block has been executed, and the status of the INDEX

## Node Description:

TOP: INDEX (pointing to the bit to be checked). INDEX $=1 \rightarrow$ The first bit.
MIDDLE: Source matrix.
BOTTOM: Matrix length (word).

## Input Control:

$\mathrm{I}_{1}$ : When $\downarrow\left(\mathrm{J}^{\top}\right)$ is presented, the function block is executed.
$I_{2}:$ INDEX control. If $I_{3}=1$ and the top node is $4 X X X X$, then the INDEX is incremented by 1 after execution.
$\mathrm{I}_{3}$ : INDEX control. $=1$, Reset INDEX.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=$ The state of the bit sensed.
$\mathrm{O}_{3}$ : Status of the INDEX.
$=1$, if INDEX is equal to zero or larger than the value of the BOTTOM node times 16 .

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 is energized and $\mathrm{I}_{2}=1$ ，the state of coil 00095 is set to that of the bit checked．Since（40421）＝0001，The 1st bit is checked．And since the bottom node is $\# 00002$ ，thus the registers $40151 \sim 40152$ are checked．


| DECO | DECODER (4->16) | $\lrcorner\llcorner$ | $\uparrow\llcorner$ |
| :---: | :---: | :---: | :---: |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| TOP | O | O | O | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE | O |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| BOTTOM |  |  |  |  | $(1)$ |  |  |

(1)0~3

## Description:

This function is a 4 bit to 16 bit decoder. The top node contains 4 sets of 4 -bit data. The set of data to be decoded is defined in the bottom node.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: Input to decoder, only four bits (nibble) are used.
MIDDLE: Decoder output.
BOTTOM: Determine which nibble in the TOP node is to be decoded.

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark(\uparrow \square)$ is presented, the function block is executed.

## Function Output:

$$
\begin{aligned}
& \mathrm{O}_{1}=\mathrm{I}_{1} \\
& \mathrm{O}_{2}=0 \\
& \mathrm{O}_{3}=0 \\
& \hline
\end{aligned}
$$

TRUTH TABLE


## 【EXAMPLE】



## 【DESCRIPTION】

Let register $40019=2$ CF9h $=0010110011111001 \mathrm{~B}$ ，and $\# 00000$ is defined in the bottom node ． \＃00000 indicates that the first set of 4－bit data is to be used as the decoder function input．The first 4－bit set in this example is 1001B，which is equal to 9 ．Therefore，the 10th bit（ 0 means the 1 st bit and 15 means the 16 th bit）in the middle node（40009）will be set after contact 10001 is energized．


These four bits are used for decoding．


## Description:

This function is a 16 -bit to 4 -bit encoder. The top node contains the data to be encoded. The bottom node indicates the 4-bit set to be used to store the encoded result, and the encoded data is stored in the middle node.

NOTE: If more than one bit is set in the top node, then the bit which is closer to the most significant bit will be used for encoding.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: Input to encoder.
MIDDLE: Encoder result.
BOTTOM: Nibble (0~3) where the encoder result is stored.

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark(\uparrow L)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=$ indicator
$=1$, if the valued stored in the TOP node is zero.
$\mathrm{O}_{3}=0$


## 【EXAMPLE】

|  | $\ddots$ |
| ---: | ---: |
| 10001 | -40009 <br> 40019 <br> ENCO <br> $\# 00000$ |

## 【DESCRIPTION】

Let $40009=0040 \mathrm{~h}=0000000001000000 \mathrm{~h}$ ，and $\# 00000$ is given in the bottom node．Since MSB is the 16th bit and LSB is the first bit in a 16－bit register，thus，the 7th bit is encoded to 6； and 6 is equal to 0110B．This 0110B 4－bit set is moved to the 1st 4－bit set of register 40019 as defined in the bottom node（\＃00000）



## SYMBOL:



OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| TOP | O | O | O | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE | O |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| BOTTOM |  |  |  |  | $(1)$ |  |  |

(1) 1~2

## Description:

This function performs binary to binary-coded-decimal conversion. The data to be converted is defined in the top node, and the converted data is stored in the middle node. The bottom node defines the conversion type (word or long word).
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed and whether the result is correct or not.

## Node Description:

TOP: data set (binary) to be converted, must be $<=9999$ (decimal).
MIDDLE: Conversion result.
BOTTOM:1.Word conversion.
2.Long word conversion.

## Input Control:

$\mathrm{I}_{1}$ : When $\sqrt[\square]{ }(\uparrow)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=$ indicator
$=1$, if the valued stored in the TOP node is $>9999$ (decimal) when the value of bottom node is ' 1 '.
$=1$, if the valued stored in the TOP node is $>99999999$ (decimal) when the value of bottom node is ' 2 '.
$\mathrm{O}_{3}=0$

## 【EXAMPLE】

| 1 | 40100 |
| :---: | :---: |
| 10001 |  |
|  | 40009 |
|  | B－＞C |
|  | \＃00002 |
|  |  |

## 【DESCRIPTION】

Let register $(40100)=0001 \mathrm{~h}$ ，and register $(40101)=0002 \mathrm{~h}$ ．When contact 10001 is energized， since \＃00002 is defined as a long word conversion，then the top node long word（40100） $10010 \mathrm{~h}=65538 \mathrm{~d}$ are converted and stored in registers 40009 and 40010.

\left.| 40100 |  |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :---: |
| 40101 | 0000 | 0000 | 0000 | 0001 |  |
|  | 0000 | 0000 | 0000 | 0010 |  |
|  |  |  |  |  |  |
|  | 0000 | 0000 | 0000 | 0110 |  |
| 40009 |  |  |  |  |  |
| 40010 | 0101 | 0101 | 0011 | 1000 |  |
|  |  |  |  |  |  |$\right\}$


|  |  | C->B |  |
| :---: | :--- | :--- | :---: |
| C->B | BCD TO BINARY CONVERTION | $\square L$ | $\uparrow L$ |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| TOP | O | O | O | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE | O |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| BOTTOM |  |  |  |  | $(1)$ |  |  |

(1) 1~2

## Description:

This function performs binary-coded-decimal to binary conversion. The data to be converted is defined in the top node, and the converted data is stored in the middle node. The bottom node defines the conversion type (word or long word).
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed and whether the result is correct or not.

## Node Description:

TOP: data set (BCD) to be converted.
MIDDLE: Conversion result.
BOTTOM:1. Word
2. Long word

## Input Control:

$\mathrm{I}_{1}$ : When $\swarrow(\uparrow L)$ is presented, the function block is executed.

## Function Output:

```
O
O
    = 1, if the valued stored in the TOP node is not in BCD format.
O
```


## 【EXAMPLE】



## 【DESCRIPTION】

Science this is a long word conversion（bottom node is \＃00002）．When contact 100d is energized， the top node long word $(30009)=65538 \mathrm{~d}$ is converted to 10010 h and stored in middle node （40019），（40020）．
Let register $(30009)=8888 \mathrm{~d}=22 \mathrm{~B} 8 \mathrm{~h}$ ，and register $(30010)=7777 \mathrm{~d}=1 \mathrm{E} 61 \mathrm{~h}$ ．
When contact 10001 is energized，since \＃00002 is defined to the bottom node，then the converted BCD codes are stored in registers 40019 and 40020.

| 30009 | 0000 | 0000 | 0000 | 0110 |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 30010 | 0101 | 0101 | 0011 | 1000 |  |
| 40019 | 0000 | 0000 | 0000 | 0001 |  |
| 40020 | 0000 | 0000 | 0000 | 0010 | 10010h |


| SSEG | SEVEN-SEGMENT DECODER | $\checkmark \_$ | $\uparrow\llcorner$ |
| :---: | :---: | :---: | :---: |

## SYMBOL:



OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| TOP | O | O | O | $\bigcirc$ |  | $\bigcirc$ |  |
| MIDDLE | O |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| BOTTOM |  |  |  |  | $(1)$ |  |  |

(1)1~4

## Description:

The register for the top node is divided into four 4-bit sets of data, and each set is converted for 7 -segment display format. The bottom node defines the size (1~4) to be converted. The result is stored in the middle node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: data to be converted.
MIDDLE: conversion result.
BOTTOM: number of digits to be converted (1digit $=4$ bits).

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark(\uparrow \square)$ is presented, the function block is executed.
$\mathrm{I}_{2}$ : Leading zero display control
$=0$, normal display.
$=1$, leading zero suppressed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$
$\mathrm{O}_{3}=0$

## 【EXAMPLE】

| 1 | 40090 |
| :---: | :---: |
| 10001 |  |
|  | 40095 |
|  | SSEG |
|  | \＃00004 |

## 【DESCRIPTION】

When $\mathrm{I}_{1}=1$ ，the data contained in register 40095 is divided into four 4－bit sets and converted to 7 －segment display format．Since the bottom node is given \＃00004，thus，all four 4－bit sets are converted and stored in the middle node registers 40095 and 40096.

BOTTOM


| Display | number | a | b | c | d | e | f | g |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{0}$ | 0 | 1 | 1 | 1 | 1 | 1 | 1 | 0 |
| $\mathbf{1}$ | 1 | 0 | 1 | 1 | 0 | 0 | 0 | 0 |
| $\mathbf{2}$ | 2 | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
| $\mathbf{3}$ | 3 | 1 | 1 | 1 | 1 | 0 | 0 | 1 |
| $\mathbf{4}$ | 4 | 0 | 1 | 1 | 0 | 0 | 1 | 1 |
| $\mathbf{5}$ | 5 | 1 | 0 | 1 | 1 | 0 | 1 | 1 |
| $\mathbf{6}$ | 6 | 1 | 0 | 1 | 1 | 1 | 1 | 1 |
| $\mathbf{7}$ | 7 | 1 | 1 | 1 | 0 | 0 | 0 | 0 |
| $\mathbf{8}$ | 8 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| $\mathbf{9}$ | 9 | 1 | 1 | 1 | 1 | 0 | 1 | 1 |
| $\mathbf{A}$ | A | 1 | 1 | 1 | 0 | 1 | 1 | 1 |
| $\mathbf{b}$ | B | 0 | 0 | 1 | 1 | 1 | 1 | 1 |
| $\mathbf{C}$ | C | 1 | 0 | 0 | 1 | 1 | 1 | 0 |
| $\mathbf{d}$ | D | 0 | 1 | 1 | 1 | 1 | 0 | 1 |
| $\mathbf{E}$ | E | 1 | 0 | 0 | 1 | 1 | 1 | 1 |
| $\mathbf{E}$ | 1 | 0 | 0 | 0 | 1 | 1 | 1 |  |

* Leading zeroes are converted to 0's.

|  |  |  |  |
| :---: | :---: | :---: | :---: |
| PACK | WORD PACK/UNPACK | $\square$ | $\uparrow \square$ |

## SYMBOL:



OPERANDS:


## Description:

Depending on $I_{2}$, this function splits the contents of the top node into two bytes, and stores them in the middle node; or, takes two LOW BYTEs from the top node, concatenate to form a new 16-bit word and stores it in the middle node.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.


## Node Description:

TOP: Data to be processed.
BOTTOM: Process result.

## Input Control:

$\mathrm{I}_{1}$ : When $\left.\square L( \rfloor L\right)$ is presented, the function block is executed.
$\mathrm{I}_{2}$ : Pack/Unpack
$=0$, Unpack (splits the source data into two words and stores them in the bottom node).
$=1$, Pack (concatenate the lower bytes of two words and stores the word in the bottom node).

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$

## 【EXAMPLE 1】

## UNPACK



## 【DESCRIPTION】

When contact 10001 is energized，the content of the top node（40090）is split into two bytes which are stored in the middle node（40095 and 40096）


BEFORE
AFTER

## 【EXAMPLE 2】

## PACK



## 【DESCRIPTION】

When contact 10001 is energized，two LOW BYTEs taken from the top node（40090 and 40091）， are concatenated to form a new 16 －bit word which is stored in the middle node（40095）．

|  |  |
| :--- | :---: |
| 40090 | 0 E 0 F |
| 40091 | 2355 |
|  |  |
| 40095 | 1234 |
|  |  |



## BEFORE

AFTER

|  |  |  | $\mathrm{I}->\mathrm{F}$ |  |
| :---: | :---: | :---: | :---: | :---: |
| $\mathrm{I}->\mathrm{F}$ | INTEGER TO FLOATING POINT |  |  |  |
| CONVERSION | $\lrcorner \mathrm{L}$ | $\uparrow \mathrm{L}$ |  |  |

SYMBOL:


OPERANDS:


## Description:

This function converts an integer stored in the top node to a floating point number and stores in the registers defined in the middle node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: Data to be converted, integer (16 bits).
BOTTOM: Conversion results (32 bits).

## Input Control:

$\mathrm{I}_{1}$ : When $\sqrt{ }(\uparrow L)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

Converts the integer stored in the top node（40120）to a floating number and stores in the middle node（40130 \＆40131）．A floating point number is represented by two words：bit0～bit6 represent the exponent，bit7is the sign bit（0：positive，1：negative），and bit8～bit31 represent the fraction．


## fraction

Sign bit exponent

Formula：

$$
I=(-1)^{\mathrm{S}} \times 2^{(\mathrm{E}-64)} \times \mathrm{Fr} \quad \mathrm{~S}=\text { sign bit, } \mathrm{E}=\text { exponent }
$$



|  |  | F->I |  |
| :---: | :---: | :---: | :---: |
| F-> I | FLOATING POINT TO INTEGER <br> CONVERSION | $\checkmark \square$ | $\uparrow L$ |

SYMBOL:

| $\mathrm{I}_{1}$ | TOP | - $\mathrm{O}_{1}$ |
| :---: | :---: | :---: |
|  | F->I |  |
| $\mathrm{I}_{2}$ | BOTTOM | $\mathrm{O}_{2}$ |

OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
|  |  |  |  |  |  |  |  |

## Description:

This function converts a floating point number stored in the top node to an integer and stores in the registers defined in the middle node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: Data to be converted, floating point number (32 bits).
BOTTOM: Conversion results. Integer (16 bits).

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark(\uparrow L)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}$ : Error output
$=1$, (overflow or < 0)

## 【EXAMPLE】



## 【DESCRIPTION】

For a floating point number C000， 0042 stored in registers 40130 and 40131 respectively，the conversion returns 0003 stored in 40150.



## SYMBOL:



OPERANDS:

(1)L1~L150

## Description:

This instruction is used to instruct the program to JUMP to the other portion of the program with matched label number and EOJ instruction.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this instruction is to be executed or not.
Function outputs can be used to determine whether the instruction has been executed.

## Node Description:

TOP: Label where JUMP is intended.

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark$ ( $\uparrow$ ) is presented, the instruction is executed.

## Function Output:

$$
\begin{aligned}
& \mathrm{O}_{1}=\mathrm{I}_{1} \\
& \mathrm{O}_{2}=0
\end{aligned}
$$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 is energized，the program between JMP L00001 and EOJ L00001 is skipped． The execution continues from PAGE N ． If contact 10001 is not energized，then no JUMP action is performed．

If contact 10001 is not energized，but contact 10002 is energized，then the program between JMP L00002 and EOJ L00002 is skipped． The execution continues from PAGE M．

PAGE N


| EOJ | END OF JUMP |  |  |
| :--- | :--- | :--- | :--- |

SYMBOL:


OPERANDS:

(1)L1~L150

## Description:

This instruction is used with JMP instruction. The label numbers must be matched. Only one JMP-EOJ pair is allowed in a ladder page. The label number must not be repeated. The program between JMP and EOJ is skipped if the Input Control condition is met.

## Node Description:

TOP: Label indicating the end of JUMP.

## Input Control:

$\mathrm{I}_{1}$ : don't care.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$

|  |  |  |  |  |  |  | SR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| SR | JUMP to SUBROUTINE | $\boxed{\square}$ | $\uparrow L$ |  |  |  |  |

SYMBOL:


OPERANDS:

(1)L1~L32

## Description:

This instruction is used to call the subroutine whose label is the same as the one defined in the Top node.
Subroutine calls may be nested, but only 16 levels are allowed.
Programming requirements are: (1)SBR and RET are paired. (2)SBR instruction is behind the JSR instruction, and (3 )RET instruction is behind the SBR instruction.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this instruction is to be executed or not.
Function outputs can be used to determine whether the instruction has been executed.

## Node Description:

TOP: Label of the subroutine to be called.

## Input Control:

$\mathrm{I}_{1}$ : When $\qquad$ L ( $\uparrow$ L ) is presented, the instruction is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

The main program area and the subroutine area are separated by the EOP instruction． If the EOP instruction does not exist，then the first SBR instruction is used as a program delimiter．

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 is energized，subroutine L00001 is executed．The program control is returned to the main program when RET L00001 is encountered．The execution of the main program is terminated when the EOP instruction is encountered．
During the execution of subroutine L00001，if contact 10002 is energized，then subroutine L00010 is executed until RET L00010 is encountered．When RET L00001 is encountered，the program control returns to the main program．
If contact 10001 is not energized，then neither subroutine L00001 nor L00010 is executed．

SYMBOL:


OPERANDS:

(1)L1~L32

## Description:

This instruction is used to define the beginning of a subroutine. A matched label RET instruction is required to define the end of the subroutine. When the subroutine is called, the program control is transferred from the main program to the next page of the program where the SBR is defined.

## Node Description:

TOP: Label of the subroutine defined.

## Input Control:

$\mathrm{I}_{1}$ : Don't care.

## Function Output:

$\mathrm{O}_{1}=0$
$\mathrm{O}_{2}=0$

|  |  | RET |  |
| :---: | :---: | :---: | :---: |
| RET | RETURN FROM SUBROUTINE |  |  |

SYMBOL:


OPERANDS:

(1)L1~L32

## Description:

This instruction is used to define the end of a subroutine. The label number is defined in the top node and must be the same as the calling SBR label number.

Node Description:
TOP: Label of the subroutine.

## Input Control:

$\mathrm{I}_{1}$ : Don't care

## Function Output:

$$
\begin{aligned}
& \mathrm{O}_{1}=0 \\
& \mathrm{O}_{2}=0 \\
& \hline
\end{aligned}
$$

|  |  |  | FOR |
| :---: | :---: | :---: | :---: |
| FOR | LOOP | L |  |

## SYMBOL:



OPERANDS:

(1) 1~255
(2) L1~L64

## Description:

The program segment between the FOR and NEXT instructions with the same label number (defined in the TOP node) is repeated for a number of times (as defined in the BOTTOM node).
Loops may be nested. Maximum of 8 levels are allowed.

## Node Description:

TOP: Label of the loop.
BOTTOM: Number of repetitions.

## Input Control:

$\mathrm{I}_{1}$ : When $\square$ ( $\uparrow \quad$ ) is presented, a matched label NEXT instruction is found whose position is behind the FOR instruction. Then, this instruction is executed.

## Function Output:

$\mathrm{O}_{1}=0$
$\mathrm{O}_{2}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001is energized，the loop（L00001）is executed 10 times，then the program resumes from page $n$ ．If contact 10001 is not energized，the loop（L00001）is skipped．If both contacts 10001 and 10002 are energized，then loop L0011 is executed N times（as defined in the bottom node 40001），and loop L00001 is executed 10 times．If contact 10001 is energized while contact 10002 is not，then the loop（L00001）is executed for 10 times，while loop L0011 is skipped．

OPERANDS:

(1)L1~L64

## Description:

This instruction is used to define the end of a loop with the same label number.

## Node Description:

TOP: Label of the loop.

## Input Control:

$\mathrm{I}_{1}$ : no action.

## Function Output:

$$
\begin{aligned}
& \mathrm{O}_{1}=\mathrm{I}_{1} \\
& \mathrm{O}_{2}=0 \\
& \hline
\end{aligned}
$$



## Description:

This instruction is used to define the end of a program. All the programming behind this instruction is ignored. The program scan terminates when this instruction is encountered.

## Node Description:

## Input Control:

$\mathrm{I}_{1}$ : Don't care

## Function Output:

$\mathrm{O}_{1}=0$
$\mathrm{O}_{2}=0$

|  |  |  |  |  |  |  |  | SKIP |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| SKIP | SKIP | $\square$ | $\uparrow$ |  |  |  |  |  |  |

## SYMBOL:



OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  | O | O | ○ | ○ |  |
|  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |

## Description:

This instruction is used to control the sequence of the program execution.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this instruction is to be executed or not.
Users are recommended to have only a SKIP instruction in a ladder page.

## Node Description:

TOP: number of program pages to be skipped. If this value is equal to 0 , then the program scan is terminated.

## Input Control:

$\mathrm{I}_{1}$ : When $\square$ $(\uparrow$ ) is presented, the instruction is executed.

## Function Output:

$$
\begin{aligned}
& \mathrm{O}_{1}=0 \\
& \mathrm{O}_{2}=0 \\
& \hline
\end{aligned}
$$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 is energized，then the skip instruction is executed．
1．If the content of register 40001 is $\# 00002$ ，then the next two pages are skipped．
2．If the content of register 40001 is 0 ，then the program execution for this scan is terminated．

Let register（40001）＝00002：
$10001=$ OFF

$\left.$|  |
| :---: | | Page 1 |
| ---: |
| Page 2 |
| Skip＇OFF＇ | \right\rvert\,

$10001=\mathrm{ON}$


The SKIP instruction is at the bottom of Page 2.

| MCS | MASTER CONTROL SET | $\checkmark\llcorner$ | $\square$ |
| :---: | :---: | :---: | :---: |

SYMBOL:


OPERANDS:


## Description:

This function block is used for controlling the program flow. There must be a matched label MSE (Master control end) function block for the ladder program to execute correctly. The power rail input of the ladder program segment between the MCS-MSE pair is determined by the $\mathrm{I}_{1}$ of MCS. If $\mathrm{I}_{1}$ is ON , the power rail input of the ladder program segment between the MCS-MSE pair is ON or rice versa. Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not. Function outputs can be used to determine whether the function block has been executed. Nesting MCS are not supported.

## Node Description:

## Input Control:

$\mathrm{I}_{1}$ : Power control

## Function Output:

$\mathrm{O}_{1}: \mathrm{I}_{1}$
$\mathrm{O}_{2}=0$


## 【DESCRIPTION】

When contact 10001 is energized, the power rail input of ladder program segment between the MCS-MSE function blocks is OFF. If the contact 10001 is not energized, then the ladder program segment is executed as usual.


SYMBOL:


OPERANDS:


## Description:

This function block is the matched ending instruction for MCS function block.

## Node Description:

Input Control:

Function Output:
$\mathrm{O}_{1}=0$
$\mathrm{O}_{2}=0$

| INIP | INITIALIZATION OF POINTER | $\downarrow \square$ | $\uparrow\llcorner$ |
| :---: | :---: | :--- | :--- |

SYMBOL:


OPERANDS:

(1)P0~P15

## Description:

This function is used to define the content of a pointer. The constant in the bottom node is used to define which pointer is to be initialized, and the number in the top node is the initialization value.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.

## Node Description:

TOP: Type of register and its number.
BOTTOM: Pointer to be defined.

## Input Control:

$\mathrm{I}_{1}$ : When $\square(\uparrow L)$ is presented, the instruction is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 is energized，the relationship：$(\mathrm{P} 0002)=40001$ is defined．
This means the（ P 0002 ）pointer points to this 40001 register．


SYMBOL:


OPERANDS:

(1)P0~P15

## Description:

This function is used to increment the pointer by one. The constant in the top node defines which pointer is to be incremented.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.

## Node Description:

TOP: Pointer to be incremented.

## Input Control:

$\mathrm{I}_{1}$ : When $\square$ $(\uparrow \quad)$ is presented, the instruction is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=$ Error
$=1$, When this function is called, the reference number pointed by the pointer is already pointed to the last reference number of that reference type.

## 【EXAMPLE】



## 【DESCRIPTION】

Assume that pointer 2 contains 40001，when $10001=1$ ．Then，pointer $2=40002$ ，i．e． P0002＝40002 after execution．

| PADD | ADDITION OF POINTER | $\checkmark\llcorner$ | $\uparrow L$ |
| :---: | :---: | :---: | :---: |

## SYMBOL:



OPERANDS:

(1)P0~P15
-0~9999

## Description:

The content of the pointer is the top node and the constant in middle node are added and the sum is stored in the content of the pointer in the bottom node.

## Node Description:

TOP: Pointer of top node
Middle: A constant
Bottom: Pointer of buttom node

## Input Control:

$\mathrm{I}_{1}$ : When $\sqrt{ }(\uparrow L)$ is presented, the function block is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}\left(\mathrm{O}_{1}\right.$ will be ' 0 ' when $\mathrm{O}_{3}$ is ' 1 ')
$\mathrm{O}_{2}=0$
$\mathrm{O}_{3}=1$ (error output)=1, if pointer is beyond the upper limited address.

## 【EXAMPLE】



## 【DESCRIPTION】

When the contact 10001 is energized，the content of pointer P0002 is added 4 and the sum is stored to the content of pointer P0003．

| DECP | DECREMENT OF POINTER | $\lrcorner \square$ | $\uparrow\llcorner$ |
| :---: | :---: | :--- | :--- |

SYMBOL:


OPERANDS:

(1)P0~P15

## Description:

This function is used to decrement the pointer by one. The constant in the top node defines which pointer is to be decremented.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.

## Node Description:

TOP: Pointer to be incremented.

## Input Control:

$\mathrm{I}_{1}$ : When $\Omega \square(\uparrow L)$ is presented, the instruction is executed.

## Function Output:

$$
\begin{aligned}
& \mathrm{O}_{1}=\mathrm{I}_{1} \\
& \mathrm{O}_{2}=\text { Error }
\end{aligned}
$$

$=1$, when this function is called, the reference number pointed by the pointer is already pointed to the first reference number of that reference type.

## 【EXAMPLE】



## 【DESCRIPTION】

Assume that pointer 2 contains 40011，when $10001=1$ ．Then，pointer $2=40010$ ，i．e． P0002＝40010 after execution．

|  |  |  | PSUB |
| :---: | :---: | :---: | :---: |
| PSUB | SUBRATION OF POINTER |  | $\uparrow$ ¢ |

SYMBOL:


OPERANDS:


## (1)P0~P15

■0~9999

## Description:

The content of the pointer in the top node is substrated by a constant in middle node and the result is stored in the content of the pointer in the bottom node.

## Node Description:

TOP: Source pointer
MIDDLE: A constant
BOTTOM: Destination pointer

## Input Control:

$\mathrm{I}_{1}$ : When $\rfloor L(\uparrow L)$ is presented, the function block is executed.

## Function Output:

```
O
O
O
```


## 【EXAMPLE】

| P0002 |
| :---: | :---: |
| $\# 00004$ <br> PSUB <br> P0003 |

## 【DESCRIPTION】

When the contact of 10001 is energized，the content of pointer P0002 is substrated 4 and the sum is stored to the content of pointer P0003．

## SYMBOL:



OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP | O | O | O | O | $\mathbb{1}$ | $\bigcirc$ |  |
|  |  |  |  |  |  |  |  |
| BOTTOM |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |

(1)0~65535

## Description:

This function is used to define the content of a register ( 4 xxxx ) or discrete output ( 0 XXXX ). Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.

## Node Description:

TOP: Referenced (or source) register or a constant.
BOTTOM: Register or (0XXXX) to be initialized (target).

## Input Control:

$\mathrm{I}_{1}$ : When $\square(\uparrow L)$ is presented, the instruction is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 is energized，the constant $\# 00100$ is stored in register 40001，i．e． $(40001)=100$ ．

OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP | $\bigcirc$ | O | $\bigcirc$ | $\bigcirc$ | $\oplus$ | $\bigcirc$ |  |
| MIDDLE | $\bigcirc$ | $\bigcirc$ | $\bigcirc$ | $\bigcirc$ | $\bigcirc$ | $\bigcirc$ |  |
| BOTTOM |  |  |  | $\bigcirc$ | $\square$ | $\bigcirc$ |  |

## (1)0~65535

—1~2

## Description:

This function is used to compare the data in the top node and the middle node.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Outputs $\left(\mathrm{O}_{1}, \mathrm{O}_{2}, \mathrm{O}_{3}\right)$ are represented the comparing result (>,=, $\langle$ ) of top node and the middle node when this function block is executed.

## Node Description:

TOP: Top node data.
MIDDLE: Middle node data.
BOTTOM: Length to be compared (1: Word, 2: Long word)

## Input Control:

$\mathrm{I}_{1}$ : When $(\uparrow L)$ is presented, the instruction is executed.

## Function Output:

$\mathrm{O}_{1}$ : comparing result (data of top node $>$ middle node)
$\mathrm{O}_{2}$ : comparing result (data of top node $=$ middle node)
$\mathrm{O}_{3}$ : comparing result (data of top node < middle node)

## 【EXAMPLE】

| $\begin{array}{lr} -1 & \vdash \\ 00080 \end{array}$ | 40001 | （ ） |
| :---: | :---: | :---: |
|  |  | 00011 |
|  | 40001 | （ ） |
|  | RCMP | 00012 |
|  | \＃00001 | （ ） |
|  |  | 00013 |

## 【DESCRIPTION】

Assumed that the content of register $(40001)=9000(10)$ and the content of register $(40002)=$ $500(10)$ ：When the contact of 00080 is energized，the coil of 00011 will be＇ ON ＇because the content of（40001）＞（40002）．

## CHAPTER 5: SYSTEM RELATED INSTRUCTIONS

|  |  |  | DGET |
| :---: | :---: | :---: | :---: |
| DGET | GET CALENDAR DATE | $\checkmark \square$ | $\uparrow$ ¢ |

## SYMBOL:



OPERANDS:


## Description:

This function is used to obtain the system date. The result is stored in the top node (two words).
The high byte of the first word represents year; the low byte of the first word represents month; the high byte of the second word represents date in a month; and the low byte of the second word represents day in a week. All numbers are in BCD format.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.

## Node Description:

TOP: Register to store the calendar date.

## Input Control:

$\mathrm{I}_{1}$ : When $\Omega(\uparrow \mathrm{L})$ is presented, the instruction is executed.


## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 is energized，the system date is copied to registers 40001 and 40002 in BCD format．Assume that the date reads：

$$
\begin{aligned}
& (40001)=1996_{\mathrm{H}} \\
& (40002)=0918_{\mathrm{H}}
\end{aligned}
$$

Then the date is Sept．18th， 1996.

## DSET

SET CALENDAR DATE

OPERANDS:
SYMBOL:


## Description:

This function is used to set the system date. The date is stored in the top node (two words).
The high byte of the first word represents year; the low byte of the first word represents month, the high byte of the second word represents date in a month; and the low byte of the second word represents day in a week. All numbers are in BCD format.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: Date, two words.

## Input Control:

$\mathrm{I}_{1}$ : When $\uparrow$ is presented, the instruction is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 is energized，the data stored in registers 40001 and 40002 are used to set the system date．
If $(40001)=1996_{H}$
$(40002)=0918_{\mathrm{H}}$
the setting date is Sept 18th， 1996.

## SYMBOL:



OPERANDS:

(1)0~65535

## Description:

This function is used to compare the system date with the reference date stored in the top node. (For data storage format, please refer to DGET or DSET functions.) The bottom node defines the MASK for comparison.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.

## Node Description:

TOP: Reference date, two words.
BOTTOM: MASK. This word is divided into three nibbles. Each nibble corresponds to the portion of the Year, Month, and Date respectively. If a corresponding nibble is zero, then that portion is ignored during comparison.

## Input Control:

$\mathrm{I}_{1}$ : When $(\square \neg L)$ is presented, the instruction is executed.

## Function Output:

$\mathrm{O}_{1}$ : Reference date $>$ System date.
$\mathrm{O}_{2}$ : Reference date $=$ System date.
$\mathrm{O}_{3}$ : Reference date $<$ System date .

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 is energized，the reference date stored registers 40001，and a constant 0111 h is compared against the system date for any difference．
If the bottom node is represented by $\mathrm{X}_{1}, \mathrm{X}_{2}, \mathrm{X}_{3}, \mathrm{X}_{4}$（hex format）then
$X_{0}:$ reserve
$X_{1}:$ nibble mask for the Year
$X_{2}:$ nibble mask for the Month
$X_{3}:$ Date

Assume the system date is Sept．18th，1996，and

$$
\begin{aligned}
40001 & =1996_{\mathrm{H}} \\
40002 & =0917_{\mathrm{H}} \\
\text { Bottom mode } & =0111_{\mathrm{H}} \quad-(\text { Day comparison is suppressed. })
\end{aligned}
$$

Since the reference date（17）is smaller than the system date（18），thus，coil 00003 is turned＇ON＇．

Note：The date comparison is made according to the following order：
Year－＞Month－＞Date
If any difference is found while comparing the higher ranked unit，the function output is set based on the comparison result，and the rest of the data is ignored．
TGET

GET SYSTEM TIME

SYMBOL:


OPERANDS:


## Description:

This function is used to obtain the system time. The result is stored in the top node (two words).
The high byte of the first word represents day in a week; the low byte of the first word represents hour, the high byte of the second word represents minute; and the low byte of the second word represent second. All numbers are in BCD format.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.

## Node Description:

TOP: Register to store the system time.

Input Control:
$\mathrm{I}_{1}$ : When $\Omega(\uparrow \mathrm{L})$ is presented, the instruction is executed.

|  | 15 14 13 12 11 10 9 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 is energized，the system time is copied to registers 40001 and 40002 in BCD format．Assume that the time reads：

$$
\begin{aligned}
& (40001)=0212_{\mathrm{BCD}} \\
& (40002)=2345_{\mathrm{BCD}}
\end{aligned}
$$

Then the system time is Tuesday， 23 minutes 45 seconds past 12 o＇clock．

## TSET

SET SYSTEM TIME f

SYMBOL:


OPERANDS:


## Description:

This function is used to set the system time. The time is stored in the top node (two words).
The high byte of the first word represents day in a week; the low byte of the first word represents hour, the high byte of the second word represents minute; and the low byte of the second word represents second.
All numbers are in BCD format.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: Register to set the system time, 2 word format as follows :

## Input Control:

$\mathrm{I}_{1}$ : When $\uparrow L$ is presented, the instruction is executed.


## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=0$

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 is energized，the data stored in registers 40001 and 40002 are used to set the system time．
If $(40001)=0212_{\text {вСD }}$
$(40002)=2345_{\text {вСD }}$
The setting date is Tuesday， 23 minutes and 45 seconds past 12 o＇clock．

| TCMP | TIME COMPARE | $\rfloor$ | $\uparrow \square$ |
| :---: | :---: | :--- | :--- |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  |  | O | $\square$ |  |  |
| MIDDLE |  |  |  |  |  |  |  |
| BOTTOM |  |  |  |  | $\square$ |  |  |

0~65535

## Description:

This function is used to compare the system time with the reference time stored in the top node. (For data storage format, please refer to TGET or TSET functions.) The bottom node defines the MASK for comparison.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.

## Node Description:

TOP: Reference time, two words.
BOTTOM: MASK. This word is divided into four nibbles. Each nibble corresponds to the portion of the Week, Hour, Minute, Second respectively. If a corresponding nibble is zero, then that portion is ignored during comparison.

## Input Control:

$\mathrm{I}_{1}$ : When $\qquad$ ( $\uparrow$ ) is presented, the instruction is executed.

## Function Output:

$\mathrm{O}_{1}$ : Reference time $>$ System time
$\mathrm{O}_{2}$ : Reference time $=$ System time
$\mathrm{O}_{3}$ : Reference time $<$ System time

## 【EXAMPLE】



## 【DESCRIPTION】

When contact 10001 is energized，the reference time stored in registers 40001，and a constant 1110 h is compared against the system time for any difference．

If the bottom node is represented by $\mathrm{X}_{1}, \mathrm{X}_{2}, \mathrm{X}_{3}, \mathrm{X}_{4}$（hex），then
$\mathrm{X}_{1}$ ：nibble mask for the day in a WEEK
$\mathrm{X}_{2}$ ：nibble mask for the HOUR
$\mathrm{X}_{3}$ ：nibble mask for the MINUTE
$\mathrm{X}_{4}$ ：nibble mask for the SECOND

Assume the system time is Tuesday， 25 minutes 34 seconds past 7 o＇clock，and

$$
\begin{aligned}
40001 & =0208_{\mathrm{BCD}} \\
40002 & =2536_{\mathrm{BCD}} \\
\text { Bottom node } & =1110_{\mathrm{H}} \quad-(\text { Second comparison is suppressed })
\end{aligned}
$$

Since the reference hour（8）is larger than the system hour（7），thus，coil 00001 is turned＇ON＇．

Note：The date comparison is made according to the following order：
Week－＞Hour－＞Minute－＞Second．
If any difference is found while comparing the higher ranked unit，the function output is set based on the comparison result，and the rest of the data is ignored．

|  |  |  | STAT |
| :---: | :---: | :---: | :---: |
| STAT | STATUS | L | $\uparrow$ |

## SYMBOL:



OPERANDS:

(1)0~63
$\mathrm{\square} 1 \sim 64$

## Description:

This function is used to obtain the system status (configuration table)
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: Address of the system configuration table where the read action is intended.
MIDDLE: Starting address of the register where the data is stored.
BOTTOM: Number of words to be read.

## Input Control:

$\mathrm{I}_{1}$ : When $\sqrt{ }(\nmid \square)$ is presented, the instruction is executed.

## Function Output:

$\mathrm{O}_{1}=\mathrm{I}_{1}$
$\mathrm{O}_{2}=$ Read indicator
$=1$, if the read action is performed beyond the configuration table limit.
$\mathrm{O}_{3}=0$

PLC STATUS Description:

| Word Order | Description |  |
| :---: | :---: | :---: |
| $\mathbf{0 0 0}$ | Second | Reserved |
| $\mathbf{0 0 1}$ | Hour | Minute |
| $\mathbf{0 0 2}$ | Month | Day of a week |
| $\mathbf{0 0 3}$ | Year | Dummy byte |
| $\mathbf{0 0 4}$ | Maximum scan time, unit: 100 us |  |
| $\mathbf{0 0 5}$ | Minimum scan time, unit: 100 us |  |
| $\mathbf{0 0 6}$ | Current scan time, unit: 100 us |  |
| $\mathbf{0 0 7}$ | PLC Elapsed time since powered on (minute) |  |
| $\mathbf{0 0 8}$ | PLC Elapsed time since powered on(hour) |  |
| $\mathbf{0 0 9}$ | Run-time Status(1) |  |
| $\mathbf{0 1 0}$ | Self-diagnosis Status(2) |  |
| $\mathbf{0 1 1}$ | PLC link - group ID |  |
| $\mathbf{0 1 2}$ | PLC link - link flag |  |
| $\mathbf{0 1 3}$ | PLC link - real time response state |  |
| $\mathbf{0 1 4}$ | Drop used state |  |
| $\mathbf{0 1 5}$ |  |  |

## Detailed Description:

## Word - 000: Reserved

Word-001~Word - 004: System date and time.

Word-005: Maximum Scan Time, unit: 100 us

Word-006: Minimum Scan Time, unit: 100 us

Word-007: Current Scan Time, unit: 100 us

Word-008 : Elapse time since powered on (minute)

Word-009 : Elapse time since powered on (hour)

Word-010~Word - 011: Run-time and Self-diagnosis:

| 15 | 14 | 13 | 12 | 11 | 10 | 09 | 08 | 07 | 06 | 05 | 04 | 03 | 02 | 01 | 00 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

bit00: RAM checksum error
bit01: Real time timer error
bit02: Watch dog timer error
bit03: Status RAM fail
bit04: Ladder RAM fail
bit05: Remote I/O module fail
bit06: Battery low
bit07: Ladder error
bit08: I/O map error
bit09: reserve
bit12: reserve
bit13: reserve
bit14: Local I/O module fail

## b. Word - 011: Self-diagnosis Status(2)


bit00: Local I/O mismatch
bit01: Installed I/O points are over the system limitation
bit02: Remote I/O mismatch
bit04: Ladder syntax error
bit05: Rack I/O mismatched
bit06 ~ bit15: Reserved

## Word-012 : Group ID (Hexadecimal) for PLC link

## Word-013 : Link flag map of master station

Bit1 to Bit15 are corresponding to the slave station \#1 to station \#15.
$\mathbf{b i t}=$ ' $\mathbf{1}$ ' : The corresponding slave station is required to be linked.
$\mathbf{b i t}=\mathbf{~} \mathbf{0}$ ' : The corresponding slave station is not required to be linked.

## Word-014 : PLC link-linking status

Bit1 to Bit15 are corresponding to the slave station \#1 to station \#15
$\mathbf{b i t}=$ ' 1 ' : Link normal, the corresponding station is required to be linked.
$\mathbf{b i t}=$ ' $\mathbf{0}$ ' : Link abnormal, the corresponding station is not required to be linked.

## Word-015 : Communication status of the remote drops

Bit1 to Bit 15 are corresponding to the drop \#1 to drop \#15.
Bit $=$ ' 1 ' the corresponding drop has occurred communication error.


## Description:

This function is similar to a 16-contact stepping switch. The source register (defined in the TOP node) is compared against the table formed by 32 target registers ( 16 sets in total, defined in the middle node), and the corresponding bit is set to ' ON ' in the bottom node.

Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Remark: The restriction of the middle node is that the content of the first register must be hero than the content of the second register for each pair.

## Node Description:

## TOP: Target register

MIDDLE: 16 pairs of data, 32 registers.
BOTTOM: Comparison result. (Assume that the top node is represented by $\mathrm{W}_{\mathrm{b}}$ and the middle node pair is represented by $\mathrm{W}_{0}, \mathrm{~W}_{1 .}$ )
$=1$ : if $\left(\left(\mathrm{W}_{\mathrm{b}}\right)>=\left(\mathrm{W}_{0}\right)\right)$ and $\left(\left(\mathrm{W}_{\mathrm{b}}\right)<\left(\mathrm{W}_{1}\right)\right)$

## Input Control:

$\mathrm{I}_{1}$ : When $\checkmark(\uparrow \square)$ is presented, the instruction is executed.

## Function Output:

$$
\begin{aligned}
& \mathrm{O}_{1}=\mathrm{I}_{1} \\
& \mathrm{O}_{2}=0 \\
& \mathrm{O}_{3}=0
\end{aligned}
$$

## 【EXAMPLE】



## 【DESCRIPTION】

The top node of the CAM switch is coming from a T1．0 timer register．It cycles from $0 \sim 155$ ．The CAM table is defined in the middle node．It starts from 0 and has an increment of 5．After comparing the content of register 40006（for example 145）and the CAM table，this function sets the $14^{\text {th }}$ bit to on and returns 16384 to the bottom node．After 5 seconds，the $15^{\text {th }}$ bit would be ＂ON＂．After another 5 seconds，the 1 st bit would be＇ON．This behavior is very similar to that of a CAM switch．

40006

40009

41000
41001
41002
41003
41004
41005
41006
41007
41008
41009
41010
41011
41012
41013
41014
41015
41016
41017
41018
41019
41020
41021
41022
41023
41024
41025
41026
41027
41028
41029
41030
41031


| CDMR | COMMON DATA MEMORY READ |  | $\uparrow L$ |
| :---: | :---: | :---: | :---: |

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  |  |  | $\mathbb{1}$ |  |  |
| MIDDLE |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| BOTTOM |  |  |  |  | $\square$ |  |  |

## (1)Please refer to the next page for detailed description.

-1~64

## Description:

This function block is used to read the common data memory from an intelligent module. Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not. Input control ( $\mathrm{I}_{2}$ ) is used to reset this function.

## Node Description:

TOP: Location of the intelligent module, please refer to the next page.
MIDDLE:3 words, please refer to the next page.
BOTTOM: Data length.

## Input Control:

$\mathrm{I}_{1}$ : When $\uparrow L$ is presented, the function is executed.
$\mathrm{I}_{2}$ : Reset
$=1$, Reset execution status.

## Function Output:

$\mathrm{O}_{1}$ : Execution status
$=1$, Executing
$=0$, non-executing
$\mathrm{O}_{2}$ : Finishing status
$=1$, Finished
$=0$, not finished yet.
$\mathrm{O}_{3}$ : Parameter status
$=0$, Parameter OK
$=1$, Parameter error.

## TOP Node: Location of the intelligent module

The location of an intelligent module is defined via Drop.Rack.Slot. The range of Drop number is from 0 to 15 . The range of Rack number is from 1 to 4 . The Range of Slot number is from 1 to 8 .

## Middle Node: Common Data Memory



Word 1: Intelligent module CDM data address offset
Word 2: Reserved.
Word 3: Target data starting address of the register for store data after read from CDM . This address is mapping to 4 xxxx registers. For example, if word $3=00010$, then the data is stored starting from register 40010.

## 【EXAMPLE】

To read the first 10 words in PAGE0 of the common data memory for an A／D module installed at Rack 2， Slot 4，and store the result in registers 40100～40109，set the middle node as follows： $(40001)=0000,(40003)=0100$ ．


Ladder program：


## 【DESCRIPTION】

Assume that $(40001)=0000$ ，and（40003）$=100$ ；when there is an OFF $\square$ ON transition for contact 10001，the CDMR function is executed．During execution，coil 00001 is＇ON＇．After execution，coil $00001=$ coil $00003=$＇ OFF ＇，and coil $00002=$＇ON＇．

SYMBOL:


OPERANDS:

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{C}$ | $\mathbf{P}$ | $\mathbf{L}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| TOP |  |  |  |  | $\mathbb{1}$ |  |  |
| MIDDLE |  |  |  | $\bigcirc$ |  | $\bigcirc$ |  |
| BOTTOM |  |  |  |  | $\square$ |  |  |

(1)Please refer to the next page for detailed description. ㅁ1~64

## Description:

This function block is used to write the common data memory from an intelligent module.
Input control ( $\mathrm{I}_{1}$ ) is used to determine whether this function block is to be executed or not.
Input control ( $\mathrm{I}_{2}$ ) is used to reset this function.
Function outputs can be used to determine whether the function block has been executed.

## Node Description:

TOP: Location of the intelligent module, please refer to the next page.
MIDDLE: 3 words, please refer to the next page.
BOTTOM: Data length.

## Input Control:

$\mathrm{I}_{1}$ : When $\uparrow$ is presented, the function is executed.
$\mathrm{I}_{2}$ : Reset
$=1$, Reset execution status.

## Function Output:

$\mathrm{O}_{1}$ : Execution status
$=1$, Executing
$=0$, non-executing
$\mathrm{O}_{2}$ : Finishing status
$=1$, Finished
$=0$, not finished yet.
$\mathrm{O}_{3}$ : Parameter status
$=0$, Parameter OK
$=1$, Parameter error.

## TOP Node: Location of the intelligent module

The location of an intelligent module is defined via Drop.Rack.Slot. The range of Drop number is from 0 to 15 . The range of Rack number is from 1 to 4 . The range of Slot number is 1 to 8 .

## Middle Node: Common Data Memory



Word 1: Intelligent module address offset
Word 2: Reserved.
Word 3: Source data starting address of the register for copy data to CDM . This address is mapping to $4 x x x x$ registers. For example, if word $3=00010$, then the data is transferred to the intelligent module starting from register 40010.

## 【EXAMPLE】

To read the 10 words data in registers 40100～40109 and write to PAGE0 of the common data memory for an A／D module installed at Rack 2，Slot 4，set the middle node as follows：

$$
(40001)=0000,(40003)=0100 .
$$



## 【DESCRIPTION】

Assume that $(40001)=0000$ ，and $(40003)=100$ ；when there is an OFF $\square \mathrm{ON}$ transition for contact 10001，the CDMW function is executed．During execution，coil 00001 is＇ON＇．After execution，coil $00001=$ coil $00003=$＇OFF＇，and coil $00002=$＇ON＇．

| PID | PID Control | $\lrcorner L$ | $\ddagger L$ |
| :---: | :---: | :---: | :---: |

## Symbol



Operands


## Descriptions:

This function block calculate the difference between the present value and the set-point, and produce control signal to minimize the difference via PID calculation.

## Nodes:

Top: PID function parameters. Please see the next page.
Middle: Working Area and Status area for PID function. Please see the following page.
Bottom: Cycle time for PID Function, unit: 1/10 sec.

## INPUT:

$\mathrm{I}_{1}$ :Auto/Manual Mode
$=1$, Output is controlled by PID function, $=0$, Output is obtained from manual input. Error detecting is still enabled.
$\mathrm{I}_{2}$ : Bumpless transition during Manual to Auto mode switching.
$=1$, Bumpless transition enabled, $=0$, Bumpless transition disabled .
$\mathrm{I}_{3}$ : Direct/Reverse Mode
$=1$, Decrease Output as Error increases. $=0$, Increase Output as Error increases.

## OUTPUT:

$\mathrm{O}_{1}:=1$, if there is any parameter error.
$\mathrm{O}_{2}:=1$, if the present value (scaled PV) is higher than the high alarm limit.
$\mathrm{O}_{3}:=1$, if the present value (scaled PV ) is lower than the low alarm limit.

## Description:

## PID Control Loop:



## PID formula:

$$
W V(\mathbf{t})=\frac{100}{\mathrm{P}_{b}}\left[\mathrm{e}(\mathbf{t})+K_{I} \int \mathrm{e}(\mathbf{t}) \mathrm{d} \mathbf{t}+\mathrm{K}_{\mathrm{b}}\left(\frac{\mathrm{de}(\mathbf{t})}{\mathrm{dt}}\right)\right]+\text { Bias }
$$

Where:
$\mathrm{MV}(\mathrm{t})=$ Control Output
$\mathrm{Pb}=$ Proportional Band
$\mathrm{e}(\mathrm{t})=$ Error (Difference between Scaled PV and SP)
$\mathrm{K}_{\mathrm{I}}=$ Constant for Integration Term, or, reset time constant
$K_{D}=$ Constant for Derivative Term, or, rate time constant
Bias $=$ Correction Value, or offset to Output

TOP Node: Register: 4xxx1~4xx16

4xxx1: An internal register used to store the scaled PV in Engineering Unit.

Scaled PV $=\frac{\text { Raw PV }}{\text { Sensor Range }} *($ UER - LER $)+$ LER
Where: Raw PV: Obtained from the difference of Register 4xx14 and 4xx16.
UER : Upper bound of Engineering measurement Range
(See also Register 4xx12)

LER : Lower bound of Engineering measurement Range (See also Register 4xx13)

Sensor Range: 4096. Assuming that an AD020 module is used to convert Raw PV signal ( $0 \sim 10 \mathrm{~V}$ ) to digital data ( $0 \sim 65535$ ), then the Raw PV must be divided by $\mathbf{1 6}$ first to maintain consistency.

4xxx2: Set Point in Engineering Unit. (0~9999)

4xxx3: PID control output MV ( $0 \sim 4096$ ). Please use a proper scaling factor to scale this control output and then send to the Output Device.
In Auto Mode ( $\mathrm{I}_{1}=1$ ), the data in this register is the result of PID calculation. In Manual Mode ( $\mathrm{I}_{1}=0$ ), filling this register by user is required.

4xxx4: High alarm limit in Engineering Unit (0~9999). This number should be greater than the Set Point.

4xxx5: Low alarm limit in Engineering Unit (0~9999). This number should be less than the Set Point.

4xxx6: Proportional Band ( $\mathrm{Pb}: 5 \sim 500$ ). The term Proportional Band is also referred to as the "sensitivity". The reciprocal of Pb is "Gain". As seen from the PID formula, the "Gain" is the proportional factor between "Error" and output MV. For example: if $\mathrm{Pb}=5$, then MV is amplified 20 times.

4xxx7: Constant for Integration Term, or, Reset time Constant ( $\mathrm{K}_{\mathrm{I}}$ : 0~9999). As seen from the PID formula, the $K_{I}$ represents the contribution of the Integral. If $K_{I}=0$, then this function block becomes a PD function block.

4xxx8: Constant for Derivative Term, or, Rate time Constant ( $\mathrm{K}_{\mathrm{D}}: 0 \sim 9999$ ). As seen from the PID formula, the $K_{D}$ represents the contribution of the Derivative. If $K_{D}=0$, then this function block becomes a PI function block. If both $K_{I}=0$ and $K_{D}=0$, then this function block becomes a proportional control function block.

4xxx9: Bias, Correction Value, or offset to Output (0~4095).

4xx10: High integral wind-up limit, or, upper bound of output. Usually this value is set at 4095.

4xx11: Low integral wind-up limit, or, lower bound of output. Usually this value is set at 0000 .

4xx12: Upper bound of Engineering Range (0001~9999). Specify the upper limit of the sensor output in Engineering Unit in this register. For example, a RTD10 module produce unsigned digital data $1500 \sim 7500$ for temperature $0^{\circ} \mathrm{C} \sim 600^{\circ} \mathrm{C}$, then specify 600 for this register. This number should
be greater than the Set Point.

4xx13: Lower bound of Engineering Range (0000~9998). Specify the lower limit of the sensor output in Engineering Unit in this register. For example, a RTD10 module produce unsigned digital data $1500 \sim 7500$ for temperature $0^{\circ} \mathrm{C} \sim 600^{\circ} \mathrm{C}$, then specify 0000 for this register. This number should be less than the Set Point.

4xx14: Raw PV. Move the data from the output PV sensor to this register. (See also Register 4xx16)

4xx15: Internal Register for storing the status of "Auto" or "Manual" mode. If the content of this register is 11 (Hex), the PID function block is in Manual mode. If 55(Hex), Auto mode.

4xx16: Correction value for Row PV. (0~4096). Specify a correction value in this register. This value is subtracted from the Raw PV (obtained from Register 4xx14), and the result is then used in the calculation of Register 4xxx1.

Middle Node: Register 4yyy1~4yyyy5

4yyy1: PID function Block Status.

Bit $1:=1$, if there is any parameter error.
Bit 2: $=1$, if High Alarm limit is exceeded.
Bit 3: $=1$, if Low Alarm limit is exceeded.
Bit 4 ~ Bit 5: Reserved.
Bit 6: =1, if PID function Block is in "Auto" mode and computing.
Bit 7 ~ Bit 12: Reserved.
Bit 13: $=\mathrm{I}_{3}$
Bit 14: $=\mathrm{I}_{2}$
Bit 15: $=\mathrm{I}_{1}$
Bit 16: Reserved.

4yyy2: Internal Register for PID Loop timer.

4yyy3: Internal Register for storing High order integral summation.

4yyy4: Internal Register for storing Low order integral summation.

4yyy5: Internal Register for storing Scaled PV used in the previous scan.

Bottom Node: Cycle time, unit: 1/10 sec. 00010 stands for one second.

## Example:





[^0]:    $\mathrm{O}_{1}=\mathrm{I}_{1}$
    $\mathrm{O}_{2}=0$
    $\mathrm{O}_{3}=0$

